



Symmetry Principles in Quantum Simulation

with Applications to the Control of Closed and Open Systems

- I. Q-Control
- II. Q-Simulation
- III. Algorithms
- IV. Applications
- Outlook
- Conclusions
- A1
- A2
- A3

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Technical University of Munich (TUM)





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Symmetry guidelines for answering:

- 1** when is a quantum hardware **universal**?
interplay of controls and coupling architecture
- 2** when can quantum system A **simulate** system B ?
in particular: **least** state-space overhead
- 3** what are the reachable sets under **collective controls**?



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Quantum Simulation

Fundamental Questions

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More generally:

what are the reachable sets in **open systems**?

Controlling Quantum Hardware

Graph Representation

I. Q-Control

Basics

Theory Made Easy

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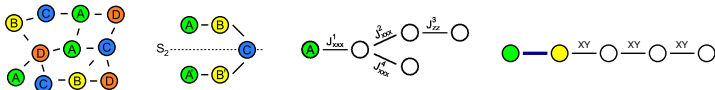
A1

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Hamiltonian components: $H_{\text{tot}} = H_0 + \sum_j u_j H_j$ in $\dot{\rho} = -i[H_{\text{tot}}, \rho]$

- vertices: **controls** = pulses (type-wise joint local actions)
- edges: **drift** = couplings (Ising-ZZ; Heisenberg-XX, XY, XXX)
- system algebra $\mathfrak{k} := \langle iH_0, iH_j \mid j = 1, 2, \dots, m \rangle_{\text{Lie}}$



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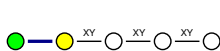
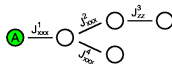
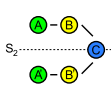
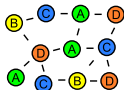
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- Consider closed control system with $\mathfrak{k} = \langle iH_\nu \rangle_{Lie}$

Definition

The *symmetry* of the Hamiltonians $\{iH_\nu\}$ is expressed by the *centraliser* (or *commutant*) of \mathfrak{k} in $\mathfrak{su}(N)$

$$\mathfrak{k}' := \{s \in \mathfrak{su}(N) \mid [s, H_\nu] = 0 \quad \forall \nu = d; 1, 2, \dots, m\}.$$

It collects all *constants of motion* under $\mathbf{K} = \langle \exp \mathfrak{k} \rangle$.



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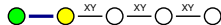
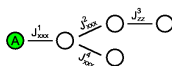
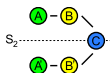
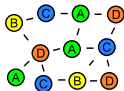
A3

Control system Σ with algebra $\mathfrak{k} = \langle iH_\nu \mid \nu = d; 1, 2, \dots, m \rangle_{\text{Lie}}$.

1 coupling graph to H_d connected

2 no symmetry (\mathfrak{k}' trivial)

(1) and (2) \Rightarrow dynamic algebra \mathfrak{k} is *simple*



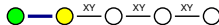
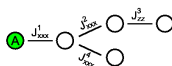
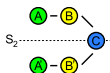
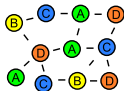


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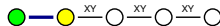
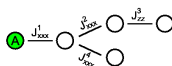
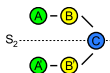
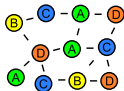


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Irreducible Simple Subalgebras to $\mathfrak{su}(N)$

up to $N = 2^{15}$

Proc. 19th MTNS Budapest 2010, 2341

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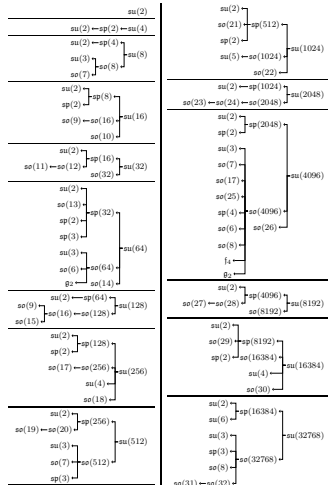
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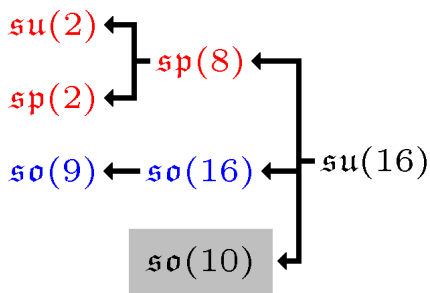
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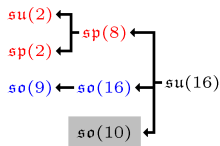
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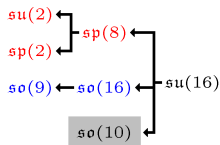
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Algorithm : Check for *conjugation* to $\mathfrak{so}(N)$ or $\mathfrak{sp}(\frac{N}{2})$
for n -qubit drift and control Hamiltonians $\{iH_d; H_1, \dots, H_m\}$

1. For each Hamiltonian $H_\nu \in \{H_d; H_1, \dots, H_m\}$
determine all non-singular *solutions to the homogeneous linear eqn.*

$$\mathcal{S}_\nu := \{\mathbf{S} \in SL(N) \mid \mathbf{S}H_\nu + H_\nu^t \mathbf{S} = 0\} \hat{=} \ker(H_\nu \otimes \mathbf{1} + \mathbf{1} \otimes H_\nu)$$

2. Check intersection of all sets of solutions

$$\mathcal{S} = \bigcap_\nu \mathcal{S}_\nu.$$

$$\text{if } \mathbf{S}\bar{\mathbf{S}} = +\mathbf{1}: \mathfrak{k} \subseteq \mathfrak{so}(N)$$

$$\text{if } \mathbf{S}\bar{\mathbf{S}} = -\mathbf{1}: \mathfrak{k} \subseteq \mathfrak{sp}(\frac{N}{2})$$

$$\text{if } \mathcal{S} = \{\} : \mathfrak{k} \text{ of other type}$$

Complexity $\mathcal{O}(N^6)$, as in Liouville space N^2 equations
have to be solved by *LU* decomposition.



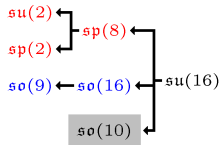
Controllability Made Easy

Necessary and Sufficient Conditions

arXiv: 1012.5256

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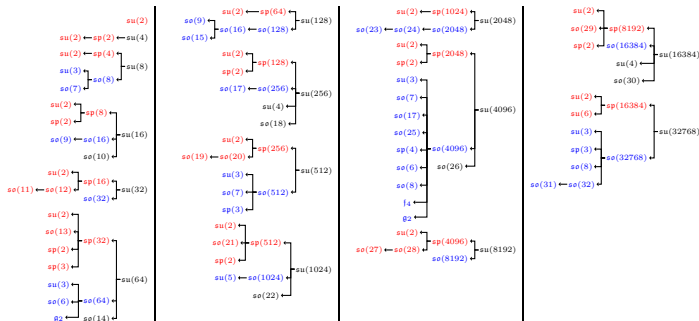
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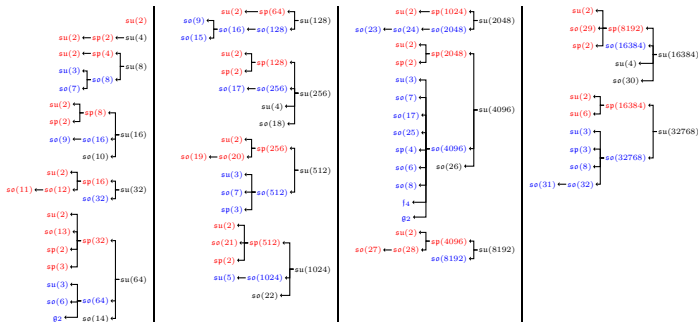
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Theorem

Let $\{H_\nu \mid \nu = d; 1, 2, \dots, m\}$ be drift and control Hamiltonians of control system Σ with system algebra \mathfrak{k} .

Define $\Phi_{AB} := \{(iH_\nu \otimes \mathbb{1}_A + \mathbb{1}_B \otimes iH_\nu) \mid \nu = d, 1, \dots, m\}$.

Then Σ is fully controllable, i.e. $\mathfrak{k} = \mathfrak{su}(2^n)$, iff

- joint commutant to Φ_{AB} is two-dimensional

i.e. $\Phi'_{AB} = \{\lambda \mathbb{1}, \text{SWAP}_{AB}\}$.

$[\Phi_{AB}] = [\text{symmetric}]_{\text{bosonic}} \oplus [\text{anti-symmetric}]_{\text{fermionic}}$



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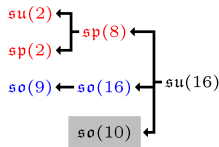
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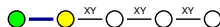
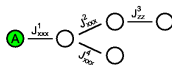
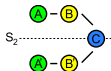
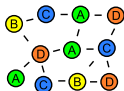
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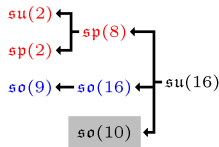
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Let Σ_A, Σ_B be control systems with *irreducible* system algebras $\mathfrak{k}_A, \mathfrak{k}_B$ over a *given Hilbert space* \mathcal{H} . Then

- 1 Σ_A simulates $\Sigma_B \Leftrightarrow \mathfrak{k}_B$ is a subalgebra of \mathfrak{k}_A ,
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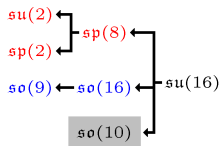
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Quantum Simulation

Overview: Local Controls

arXiv: 1012.5256

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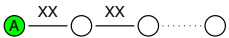
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system type	no. of levels	fermionic	bosonic	system alg.
n -spins- $\frac{1}{2}$		order of coupling		
	n	quadratic (i.e. 2)	—	$so(2n + 1)$
	$n + 1$	quadratic (i.e. 2)	—	$so(2n + 2)$
for $n \bmod 4 \in \{0, 1\}$	n	up to n	—	$so(2^n)$
for $n \bmod 4 \in \{2, 3\}$	n	—	up to n	$sp(2^{n-1})$
	n	up to n	up to n	$su(2^n)$

NB: no. of spins maps into no. of levels (as in Jordan-Wigner transformation).



Quantum Simulation

Overview: Local Controls

arXiv: 1012.5256

I. Q-Control

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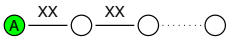
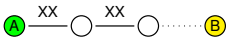
Outlook

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system type		no. of levels	fermionic order of coupling	bosonic	system alg.
n -spins- $\frac{1}{2}$					
		n	quadratic (i.e. 2)	—	$so(2n+1)$
		$n+1$	quadratic (i.e. 2)	—	$so(2n+2)$
	for $n \bmod 4 \in \{0, 1\}$	n	up to n	—	$so(2^n)$
	for $n \bmod 4 \in \{2, 3\}$	n	—	up to n	$sp(2^{n-1})$
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system type		no. of levels	fermionic ——— order of coupling ———	bosonic	system alg.
n -spins- $\frac{1}{2}$					
	XX	n	quadratic (i.e. 2)	—	$so(2n+1)$
	XX	$n+1$	quadratic (i.e. 2)	—	$so(2n+2)$
	XX				
for $n \bmod 4 \in \{0, 1\}$		n	up to n	—	$so(2^n)$
for $n \bmod 4 \in \{2, 3\}$		n	—	up to n	$sp(2^{n-1})$
		n	up to n	up to n	$su(2^n)$

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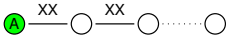
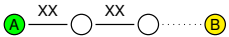
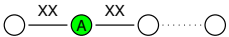
Outlook

Conclusions

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system type		no. of levels	fermionic —— order of coupling ——	bosonic	system alg.
n -spins- $\frac{1}{2}$		n	quadratic (i.e. 2)	—	$so(2n+1)$
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		n	up to n	—	$so(2^n)$
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system type		fermionic	bosonic	system alg.
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	n	quadratic (i.e. 2)	—	$so(2n+1)$
	$n+1$	quadratic (i.e. 2)	—	$so(2n+2)$
	for $n \bmod 4 \in \{0, 1\}$	up to n	—	$so(2^n)$
	for $n \bmod 4 \in \{2, 3\}$	—	up to n	$sp(2^n-1)$
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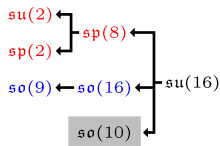
A1

A2

A3

system type	no. of levels	fermionic order of coupling	bosonic	system alg.
n -spins- $\frac{1}{2}$		—	—	
	n	quadratic (i.e. 2)	—	$so(2n+1)$
	$n+1$	quadratic (i.e. 2)	—	$so(2n+2)$
	n	up to n	—	$so(2^n)$
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for $n \bmod 4 \in \{2, 3\}$	n	—	up to n	$su(2^n)$
	n	up to n	up to n	$su(2^n)$

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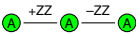


Quantum Simulation

Overview: Collective Controls for Bosonic Systems

arXiv: 1012.5256

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system type	no. of levels	bosonic coupling order	system alg. $sp(2^n - 1)$
$n = 2k + 1$ spins- $\frac{1}{2}$			
	$n = 3$	up to $n = 3$	$sp(8/2)$
	"	"	"
	$n = 5$	up to $n = 5$	$sp(32/2)$
	"	"	"
	"	"	"
	"	"	"
	"	"	"
	"	"	"

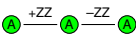
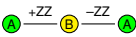


Quantum Simulation

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	"	"	"
	"	"	"
	"	"	"
	"	"	"
	"	"	"
	"	"	"

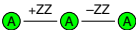
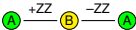
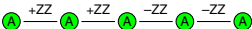


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system type	no. of levels	bosonic coupling order	system alg. $sp(2^n - 1)$
$n = 2k + 1$ spins- $\frac{1}{2}$ 	$n = 3$	up to $n = 3$	$sp(8/2)$
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	$n = 5$	up to $n = 5$	$sp(32/2)$
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"



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	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"



Quantum Simulation

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system type	no. of levels	bosonic coupling order	system alg. $sp(2^{n-1})$
$n = 2k + 1$ spins- $\frac{1}{2}$ 	$n = 3$	up to $n = 3$	$sp(8/2)$
	—"	—"	—"
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	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
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	—"	—"	—"
—"	—"	—"	—"
—"	—"	—"	—"
—"	—"	—"	—"

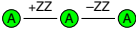

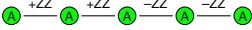
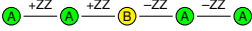
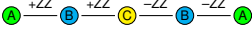
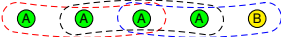
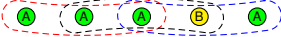


Quantum Simulation

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	—"	—"	—"
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	—"	—"	—"



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	—"	—"	—"
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	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"
	—"	—"	—"



Quantum Simulation

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	—	—	—
	—	—	—
	—	—	—
	—	—	—
	—	—	—



I. Q-Control

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Concept

Results I: conc. vs seq.

Results II: Grad. Calcs.

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Consider:

1 **linear** control system: $\dot{x}(t) = Ax(t) + Bv$

2 **bilinear** control system: $\dot{X}(t) = (A + \sum_j u_j B_j)X(t)$

Conditions for Full Controllability:

1 in **linear** systems: $\text{rank} [B, AB, A^2B, \dots, A^{N-1}B] = N$

2 in **bilinear** systems: $\langle A, B_j \mid j = 1, 2, \dots, m \rangle_{\text{Lie}} = \mathfrak{k} = \mathfrak{su}(N)$



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1 **linear** control system: $\dot{x}(t) = Ax(t) + Bv$

2 **bilinear** control system: $\dot{X}(t) = (A + \sum_j u_j B_j)X(t)$

e.g.: Ham. quantum system $\dot{U}(t) = -i(H_d + \sum_j u_j H_j)U(t)$

open quantum system $\dot{F}(t) = -(i \operatorname{ad}_{H_d} + i \sum_j u_j \operatorname{ad}_{H_j} + \Gamma_L)F(t)$

Conditions for Full Controllability:

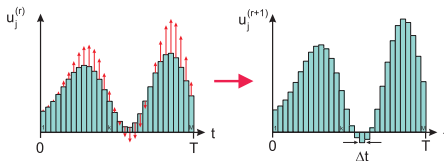
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J. Magn. Reson. **172** (2005), 296 and *Phys. Rev. A* **72** (2005), 042331

concurrent (GRAPE)



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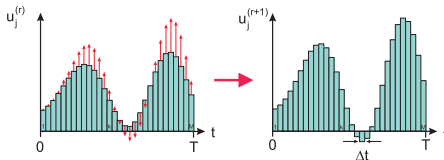
A2

A3

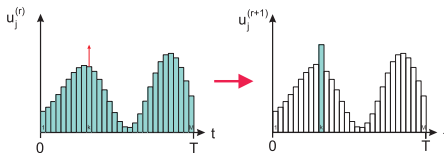


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sequential



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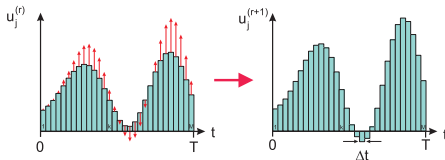
A2

A3

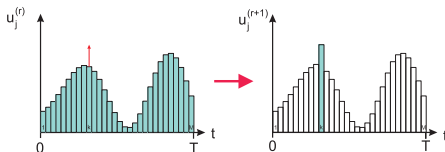


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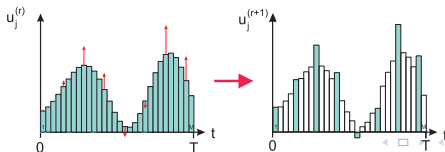
concurrent (GRAPE)



sequential



hybrid



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DYNAMO: Unified Platform

Modules for Unconstrained Bilinear Control

I. Q-Control

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0. **initialise** amplitudes $u_j^{(0)}(t_k) \in \mathcal{U} \subseteq \mathbb{R}$ for all times t_k
with $k \in \mathcal{T}_k^{(0)} := \{1, 2, \dots, M\}$, def. $X_0, X_{\text{tar}}^\dagger$.

1. **exponentiate** $X_k = e^{\Delta t A_u(t_k)}$ for all $k \in \mathcal{T}_k^{(r)}$ with
 $A_u(t_k) := A + \sum_j u_j(t_k) B_j$

2. **multiplication I** $X_{k:0} := X_k \cdot X_{k-1} \cdots X_1 (\cdot X_0 = \mathbf{I})$

3. **multiplication II** $\Lambda_{M+1:k+1}^\dagger := X_{\text{tar}}^\dagger \cdot X_M \cdot X_{M-1} \cdots X_{k+1}$

4. **evaluate fidelity** $f = \frac{1}{N} |\text{tr} \{ \Lambda_{M+1:k+1}^\dagger X_{k:0} \}|$

5. **approximate gradients** $\frac{\partial f(X(t_k))}{\partial u_j}$ for all $k \in \mathcal{T}_k^{(r)}$

6. **update amplitudes** for all $k \in \mathcal{T}_k^{(r)}$

$$\text{e.g. } u_j^{(r+1)}(t_k) = u_j^{(r)}(t_k) + F(\alpha_k, \text{Hess}_k^{-1}, \frac{\partial f(X(t_k))}{\partial u_j})$$

7. **loops**

inner: while $\| \frac{\partial f_k}{\partial u_j} \| > g_{\text{limit}}$ for $k \in \mathcal{T}_k^{(r)}$ **goto step 1** ($s \mapsto s+1$)

outer: else **goto step 1** with new set $\mathcal{T}_k^{(r+1)}(r \mapsto r+1)$



DYNAMO: Unified Platform

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DYNAMO: Unified Platform

Modules for Unconstrained Bilinear Control

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II. Q-Simulation

III. Algorithms

Concept

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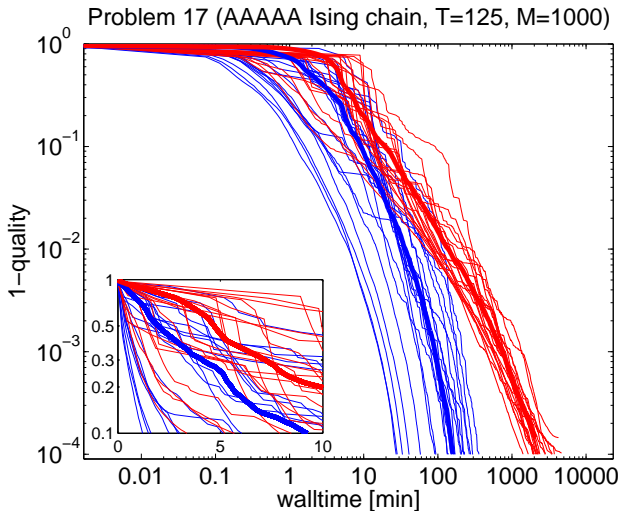


Comparison: 5-Qubit QFT

Ising Chain with Stark Shift

PRA 84 022305 (2011)

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■ concurrent (GRAPE)

■ sequential (KROTOV)



Comparison: Random Unitary

Driven Spin-6 System: model of Seth M. and Frank W.

I. Q-Control

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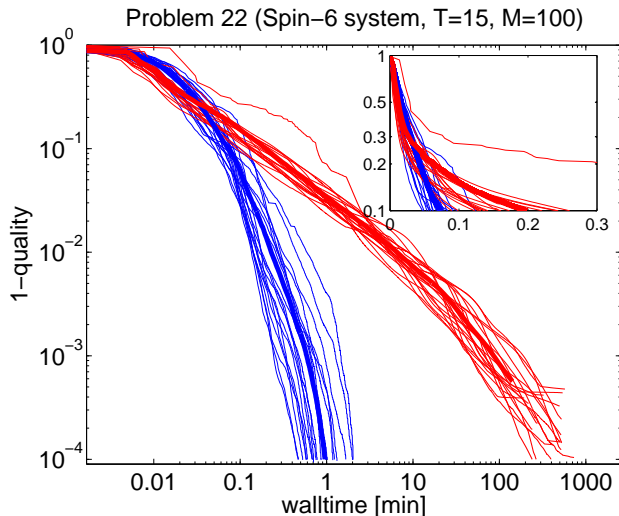
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Comparison: Random Unitary

AB000 Heisenberg-XXX Chain

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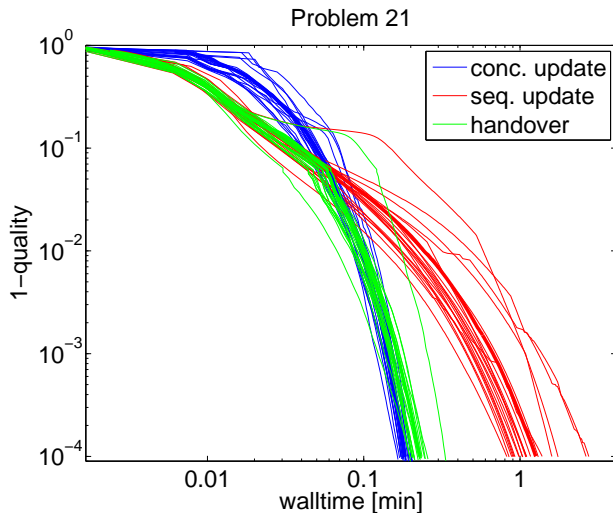
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Driven Spin-6 System

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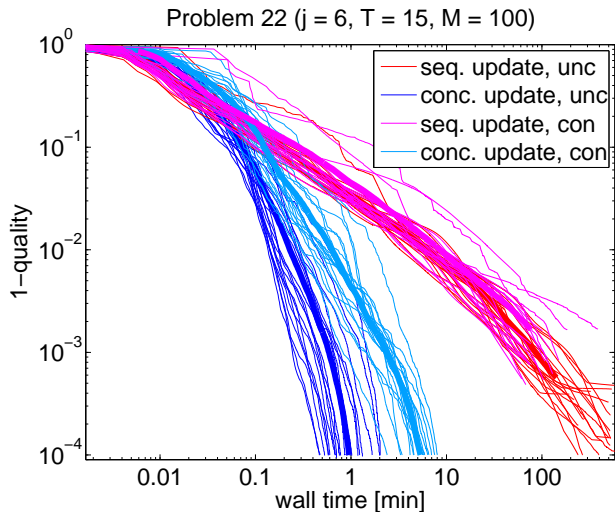
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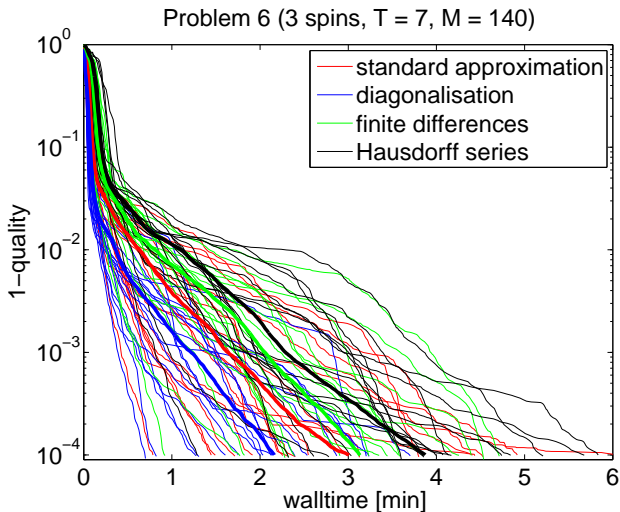


Gradient Calculation

Exact Gradients Pay

PRA 84 022305 (2011)

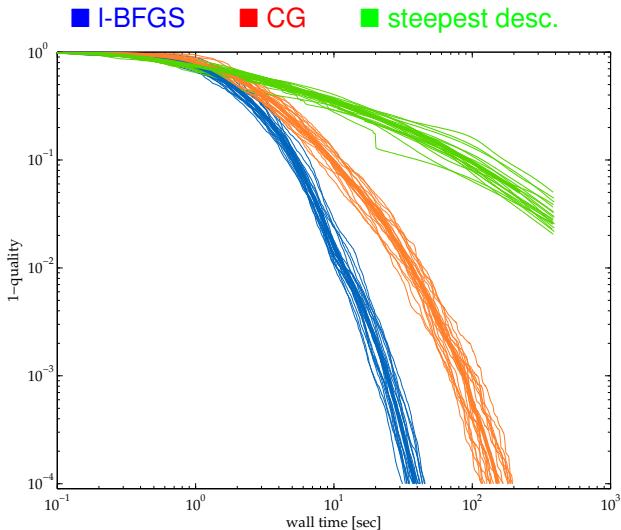
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Comparison to Conjugate Gradients

Driven Spin-3 System: Rand. Unitary Uwe Sander's PhD Thesis (2010)



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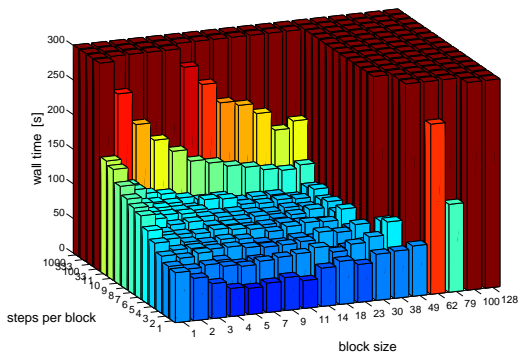
A1

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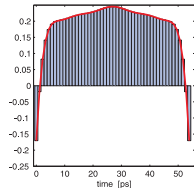
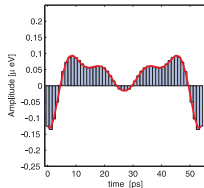
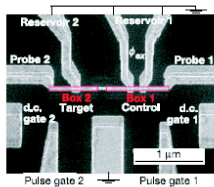
■ first-order hybrids



Time-Optimal Quantum Control

Realising Quantum Gates for Charge Qubits

■ set-up



⇒ timeopt. CNOT: **some 5 times faster** than NEC group

- Quality $q := Fe^{-\tau_{\text{op}}/\tau_Q}$

$$\text{error } 1 - q = 1 - 0.999999999 e^{-55\text{ps}/10\text{ns}} = \mathbf{0.0055}$$

$$(\text{NEC: } 1 - q = 1 - 0.4188 e^{-250\text{ps}/10\text{ns}} = 0.5917)$$

PRA 75, 012302 (2007)



Time-Optimal Quantum Control

Realising Quantum Gates for Charge Qubits

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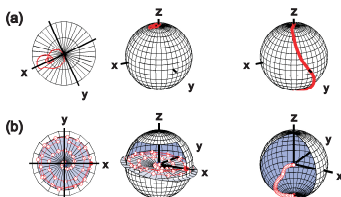
Conclusions

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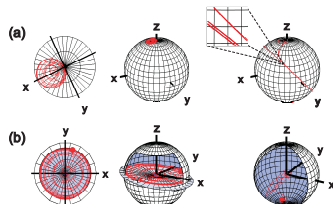
A2

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time-optimal



NEC pioneer group



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PRA 75, 012302 (2007)



Time-Optimal Quantum Control

Realising Quantum Gates for Charge Qubits with F. Wilhelm, M. Storz

Goal: realise *timeoptimal* CNOT on 2 coupled charge qubits

■ pseudospin Hamiltonian: $H = H_{\text{drift}} + H_{\text{control}}$

$$\begin{aligned} H_{\text{drift}} = & - \left(\frac{E_m}{4} + \frac{E_{c1}}{2} \right) (\sigma_z^{(1)} \otimes \mathbf{1}) - \frac{E_{J1}}{2} (\sigma_x^{(1)} \otimes \mathbf{1}) \\ & - \left(\frac{E_m}{4} + \frac{E_{c2}}{2} \right) (\mathbf{1} \otimes \sigma_z^{(2)}) - \frac{E_{J2}}{2} (\mathbf{1} \otimes \sigma_x^{(2)}) \\ & + \frac{E_m}{4} (\sigma_z^{(1)} \otimes \sigma_z^{(2)}) \end{aligned}$$

$$\begin{aligned} H_{\text{control}} = & \left(\frac{E_m}{2} n_{g2} + E_{c1} n_{g1} \right) (\sigma_z^{(1)} \otimes \mathbf{1}) \\ & + \left(\frac{E_m}{2} n_{g1} + E_{c2} n_{g2} \right) (\mathbf{1} \otimes \sigma_z^{(2)}) \end{aligned}$$

NB: components $\{H_d + H_c, H_c\}$ form minimal generating set of $\mathfrak{su}(4)$.

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Time-Optimal Quantum Control

Realising Quantum Gates for Charge Qubits

■ Symmetry: real symmetric Hamiltonians

- ⇒ palindromic controls for self-inverse gates (CNOT)
- ⇒ composed of cos Fourier series
- ⇒ Cauer synthesis by LC elements (no resistive R)

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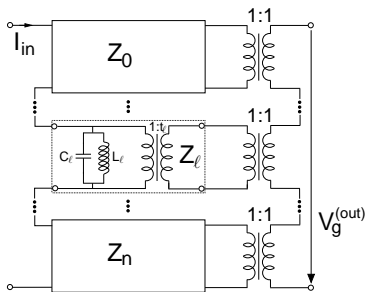


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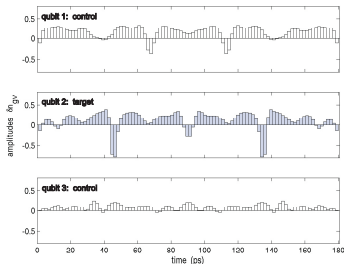
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Goal: TOFFOLI gate on 3 linearly coupled charge qubits



13 times faster than NEC

■ error rates cut by two orders of magnitude ($T_2 \simeq 10$ ns):

1 direct gate by optimal control

$$1 - q = 1 - 0.99999 e^{-180\text{ps}/10\text{ns}} = 0.0178$$

2 by 9 CNOT's from optimal control

$$1 - q = 1 - (0.999999999 e^{-55\text{ps}/10\text{ns}})^9 = 0.0483$$

3 by 9 CNOT's under pioneering controls

$$1 - q = 1 - (0.4188 e^{-250\text{ps}/10\text{ns}})^9 = 0.9997$$



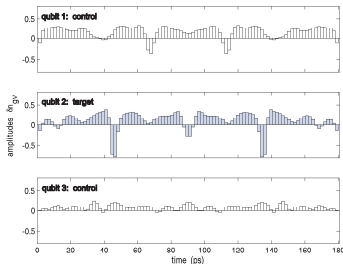
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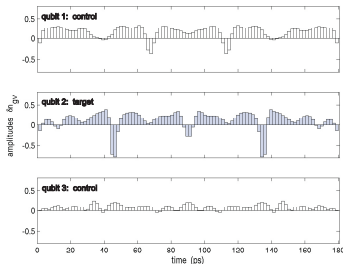
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$$1 - q = 1 - (0.999999999 e^{-55\text{ps}/10\text{ns}})^9 = 0.0483$$

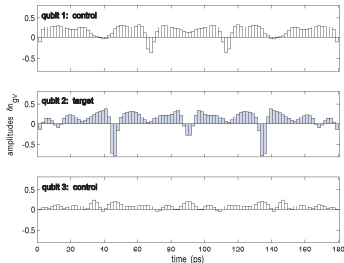
3 by 9 CNOT's under pioneering controls

$$1 - q = 1 - (0.4188 e^{-250\text{ps}/10\text{ns}})^9 = 0.9997$$

Examples of Quantum Control

Realising Quantum Gates for Charge Qubits

Goal: TOFFOLI gate on 3 linearly coupled charge qubits



13 times faster than NEC

■ error rates cut by **two orders of magnitude** ($T_2 \simeq 10$ ns):

1 direct gate by optimal control

$$1 - q = 1 - 0.99999 e^{-180\text{ps}/10\text{ns}} = \mathbf{0.0178}$$

2 by 9 CNOT's from optimal control

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Principles: Optimal Quantum Control

Scope in Optimal Control:

maximise quality function **subject to** equation of motion

Scenarios:

■ Hamiltonian dynamics

notation: $U := e^{-itH}$; $\text{Ad}_U(\cdot) := U(\cdot)U^{-1}$; $\text{ad}_H(\cdot) := [H, \cdot]$

1. pure state $|\dot{\psi}\rangle = -iH |\psi\rangle \in \mathcal{H}$
2. gate $\dot{U} = -iH U \in \mathcal{U}(\mathcal{H})$
3. non-pure state $\dot{\rho} = -i \text{ad}_H(\rho) \in \mathcal{B}_1(\mathcal{H})$
4. projective gate $\dot{\text{Ad}}_U = -i \text{ad}_H \circ \text{Ad}_U \in \mathcal{U}(\mathcal{B}_1(\mathcal{H}))$

■ Master equations of dissipative dynamics

- 3'. non-pure state $\dot{\rho} = -(i \text{ad}_H + \Gamma)(\rho)$
- 4'. **contractive** map $\dot{F} = -(i \text{ad}_H + \Gamma) \circ F \in \mathcal{GL}(\mathcal{B}_1(\mathcal{H}))$

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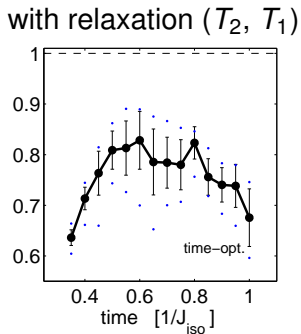
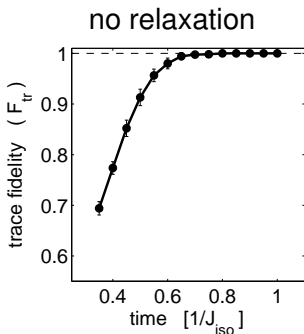
A2

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Examples of Quantum Control

3. Decoherence Control: Results of System II

■ System-II: driving **outside** slowly-relaxing subspace

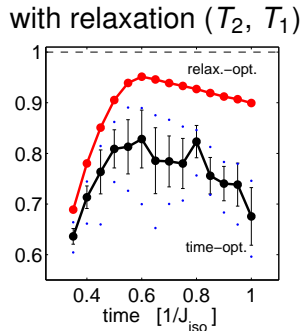
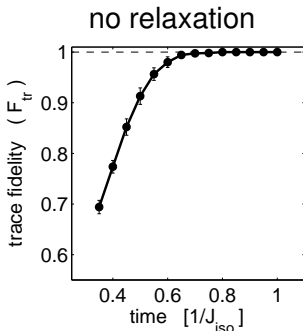


- mean of 15 time-optimised pulse sequences
- dissipation affects sequences differently

Examples of Quantum Control

3. Decoherence Control: Results of System II

■ System-II: driving **outside** slowly-relaxing subspace



- mean of 15 time-optimised pulse sequences
- dissipation affects sequences differently
- relaxation-optimised: **systematic substantial gain**

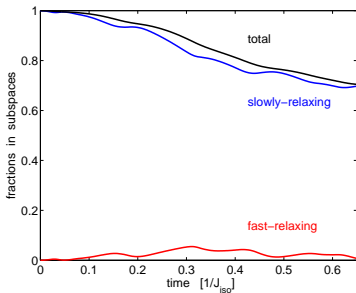


Examples of Quantum Control

3. Realising Quantum Gates with Minimal Relaxation

CNOT under **System-II**: Projection into Subspaces

■ time-optimised



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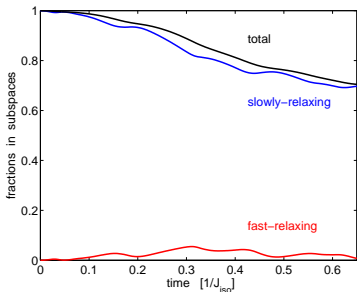
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Examples of Quantum Control

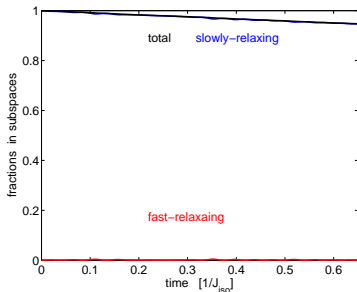
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CNOT under **System-II**: Projection into Subspaces

■ time-optimal



■ opt. against decoherence



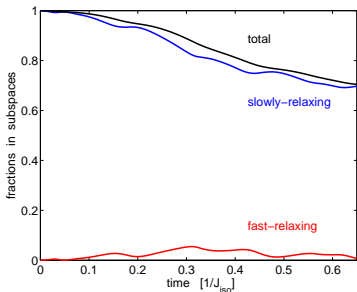
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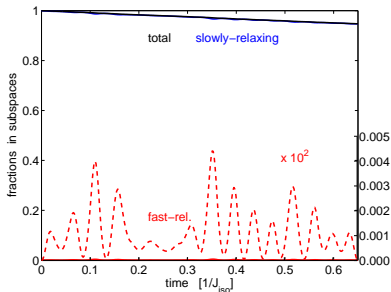
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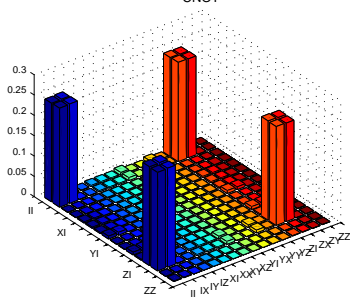
Examples of Quantum Control

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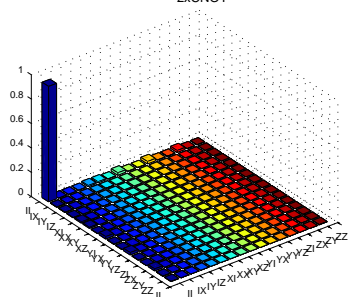
JPB 44 154013 (2011), quant-ph/0609037

■ CNOT under **System-II**: Process Tomography of Gate Protected against Dissipation by Optimal Control

$\text{abs}(\chi_{\text{CNOT}})$



$\text{abs}(\chi_{2\times\text{CNOT}})$



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Examples of Quantum Control

3. Realising Quantum Gates with Minimal Relaxation

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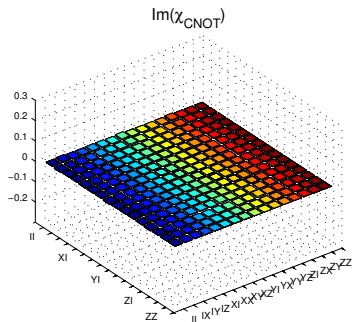
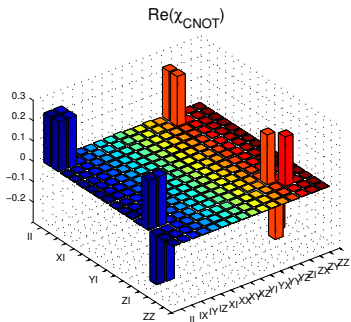
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Examples of Quantum Control

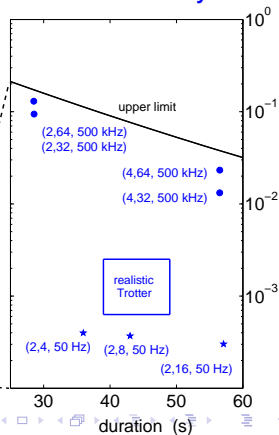
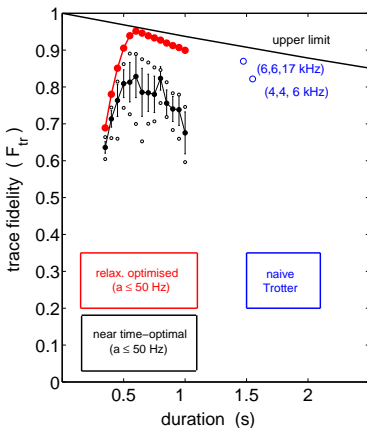
3. Realising Quantum Gates with Minimal Relaxation

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■ CNOT under **System-II**: comparison of methods

by decoherence control:
> 95% fidelity

conventional:
< 15% fidelity



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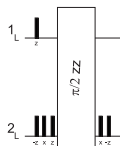


Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ logical qubits



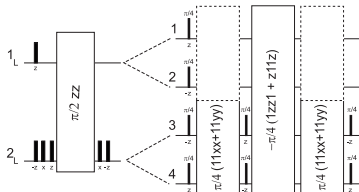
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Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ physical qubits



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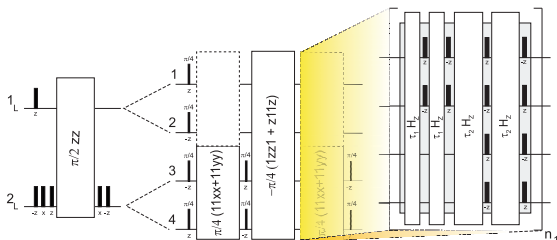
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Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ realisation by **System-I**



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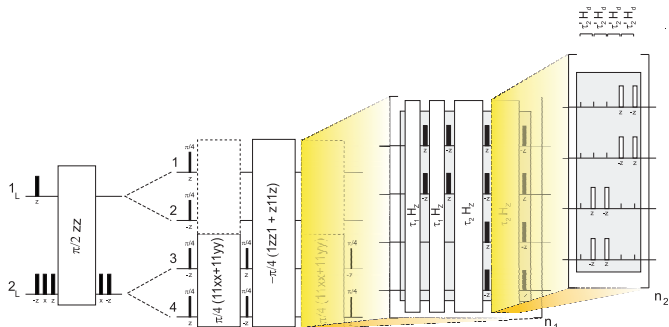


Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

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■ realisation by **System-II**



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Control of Non-Markovian Open Systems

Qubit Coupled via Two-Level Fluctuator to Spin Bath

with P. Rebentrost and F. Wilhelm

Model:

qubit coupled to a two-level fluctuator coupled to a bath

$$H = H_S + H_I + H_B$$

- $H_S = E_1(t)\sigma_Z + \Delta\sigma_X + E_2\tau_Z + \Lambda\sigma_Z\tau_Z$
- $H_I = \sum_i \lambda_i (\tau^+ b_i + \tau^- b_i^\dagger)$
- $H_B = \sum_i \hbar\omega_i b_i^\dagger b_i$

Ohmic bath spectrum: $J(\omega) = \sum_i \lambda_i^2 \delta(\omega - \omega_i) = \kappa\omega\Theta(\omega - \omega_c)$

couplings λ_i , damping κ , high-freq. cut-off ω_c

PRL 102 090401 (2009)

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Control of Non-Markovian Open Systems

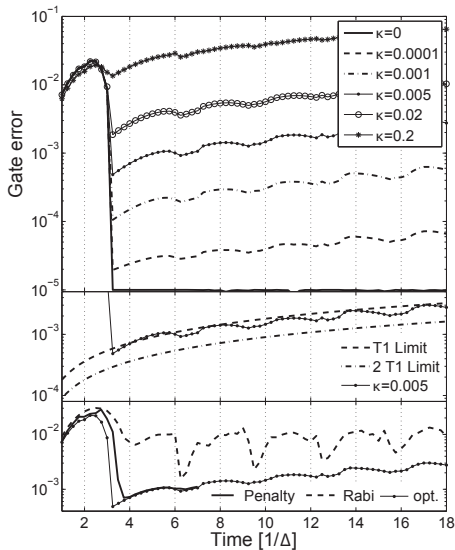
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← RABI pulse
← cut error by factor ≤ 10
with optimal control

Control of Non-Markovian Open Systems

- Principle: embed to **Markovian** and **project**

$$\begin{array}{ccc}
 \rho_0 = \rho_{SE}(0) \otimes \rho_B(0) & \xrightarrow{\text{Ad}_W(t)} & \rho(t) = W(t)\rho_0 W^\dagger(t) \\
 \Pi_{SE} \downarrow \text{tr}_B & & \Pi_{SE} \downarrow \text{tr}_B \\
 \rho_{SE}(0) & \xrightarrow[\text{Markovian}]{F_{SE}(t)} & \rho_{SE}(t) \\
 \Pi_S \downarrow \text{tr}_E & & \Pi_S \downarrow \text{tr}_E \\
 \rho_S(0) & \xrightarrow[\text{non-Markovian}]{F_S(t)} & \rho_S(t)
 \end{array}$$

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Consider: controlled system with *time dep* Lindbladians $\{\mathcal{L}_u(t)\}$

$$\dot{X} = -\mathcal{L}_u(t)X = -(iH_d + i \sum_j u_j(t)H_j + \Gamma)X$$

Lindbladians $\{\mathcal{L}_u\}$ form

- *Lie wedge* \mathfrak{w}
- *Lie semialgebra* \mathfrak{w}_s , if $\{\mathcal{L}_u\}$ BCH compatible with \mathfrak{w}
i.e. $L_j * L_k := L_j + L_k + \frac{1}{2}[L_j, L_k] + \dots \in \mathfrak{w}$
then $\{e^{-t\mathcal{L}_{\text{eff}}} \mid t > 0\}$ *physical* at all times.
- Else $\{e^{-t\mathcal{L}_{\text{eff}}} \mid t > 0\}$ *unphysical* except $t = 0$; $t = t_{\text{eff}}$ etc.



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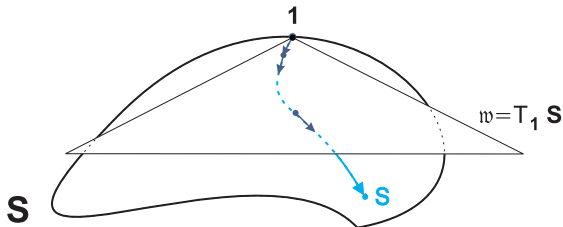


Lindbladians Generate (Lie) Semigroup

Rep. Math. Phys. 64 93 (2009)

Consider: controlled system with *time dep* Lindbladians $\{\mathcal{L}_u(t)\}$

$$\dot{X} = -\mathcal{L}_u(t)X = -(iH_d + i\sum_j u_j(t)H_j + \Gamma)X$$



Lindbladians $\{\mathcal{L}_u\}$ form

■ *Lie wedge* \mathfrak{w}

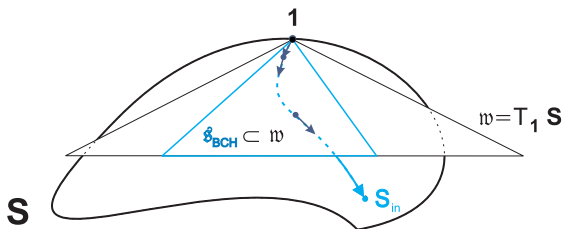
■ *Lie semialgebra* \mathfrak{w}_S

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Consider: controlled system with *time dep* Lindbladians $\{\mathcal{L}_u(t)\}$

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Lindbladians $\{\mathcal{L}_u\}$ form

- *Lie wedge* w
- *Lie semialgebra* w_S

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Lie – Markov Correspondence

Quantum Channels as Lie Semigroups

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Corollary (*cave*: ‘woodcut’, details in Rep. Math. Phys. **64** (2009) 93.)

- A channel is (time dependent) **Markovian**, iff there is representation $T = e^{-\mathcal{L}_1} e^{-\mathcal{L}_2} \dots e^{-\mathcal{L}_r}$ so that the $\mathcal{L}_1, \mathcal{L}_2, \dots, \mathcal{L}_r$ generate a **Lie wedge** \mathfrak{w}_r .
- Moreover, T specialises to *time independent* form, iff its **Lie wedge** \mathfrak{w}_r specialises to a *Lie semialgebra*.

Complements recent work: Wolf,Cirac, *Commun. Math. Phys.* (2008) & Wolf,Eisert,Cubitt,Cirac, *PRL* (2008)



Lie – Markov Correspondence

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Consider bilinear control system

$$\dot{X} = -(A + \sum_j u_j B_j)X \text{ with } A := i\hat{H}_d + \Gamma_L \text{ and } B_j := i\hat{H}_j$$

- controllability condition for **closed** systems:

$$\langle iH_d, iH_j \mid j = 1, 2, \dots, m \rangle_{\text{Lie}} = \mathfrak{su}(N)$$

- **WH**-condition for **open** systems:

$$\langle iH_d, iH_j \mid j = 1, 2, \dots, m \rangle_{\text{Lie}} = \mathfrak{su}(N)$$

- **H**-condition for **open** systems:

$$\langle iH_j \mid j = 1, 2, \dots, m \rangle_{\text{Lie}} = \mathfrak{su}(N)$$



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- closed controllable systems:

$$\text{Reach } \rho_0 = \mathcal{O}_U(\rho_0) := \{U\rho_0 U^\dagger \mid U \in SU(N)\}$$

- open fully H-controllable systems:

$$\text{Reach } \rho_0 \subseteq \{\rho \in \text{pos}_1 \mid \rho \prec \rho_0\}$$

- open systems satisfying WH-condition:
parameterisation involved, key: Lie semigroups



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Exploring Reachable Sets

by Lie Semigroups

arXiv: 1103.2703

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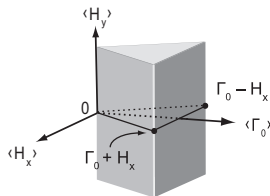
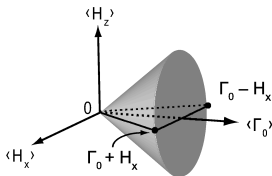
Bilinear control system: $\dot{X} = -(A + \sum_j u_j B_j)X$

- satisfies WH-condition with :

$$A := H_z + \Gamma_0, B := uH_y, \text{ and } \Gamma_0 := \text{diag}(1, 0, 1)$$

- Lie wedge:

$$\mathfrak{w}_0 = \langle H_y \rangle \oplus -\mathbb{R}_0^+ \text{conv} \left\{ \begin{bmatrix} \sin(\theta) \\ \cos(\theta) \\ 1 \end{bmatrix} \cdot \begin{bmatrix} H_x \\ H_z \\ \Gamma_0 \end{bmatrix} \mid \theta \in \mathbb{R} \right\}$$





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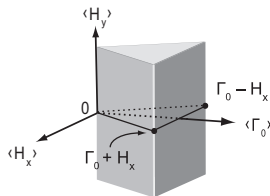
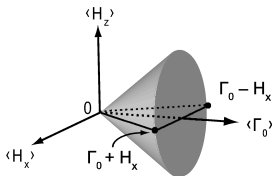
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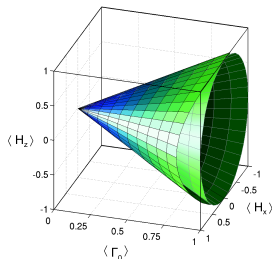
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$$A := H_z + \Gamma_0, B := uH_y, \text{ and } \Gamma_0 := \text{diag}(1, 1, 2)$$

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$$\mathfrak{w}_0 = \langle H_y \rangle \oplus -\mathbb{R}_0^+ \text{conv} \left\{ \begin{bmatrix} 2 \sin(\theta) \\ 2 \cos(\theta) \\ \gamma \sin(2\theta) \\ \gamma(1 - \cos(2\theta)) \\ (11 + \cos(2\theta))/6 \end{bmatrix} \cdot \begin{bmatrix} H_x \\ H_z \\ \rho_y \\ \Delta \\ \Gamma_0 \end{bmatrix} \mid \theta \in \mathbb{R} \right\}$$



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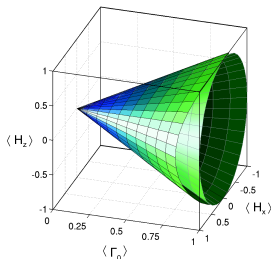
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$$\mathbf{S} = e^{A_1} e^{A_2} \dots e^{A_\ell} \text{ with } A_1, A_2, \dots, A_\ell \in \mathfrak{w}$$



Conclusions

Exploit *symmetries* for the Royal Road to:

1 Controllability

- absence of symmetry plus inclusion **fermionic** & **bosonic** systems

2 Simulability

- efficient q-simulation: algebraic understanding

3 DYNAMO Modular Platform

4 Lie Semigroups

- as new *unifying framework* for **open** systems

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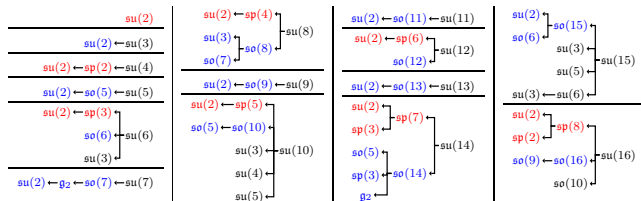
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Gunther Dirr, Corey O'Meara

integrated EU programme; excellence network; high-speed parallel cluster



References:

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arXiv:0904.4654, IEEE Proc. ISCCSP 2010 23.2, Proc. MTNS, 2341 (2010),

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Design Rules

For an n spin- $\frac{1}{2}$ system with a **connected coupling graph** and **no symmetries** to be fully controllable it suffices that

- (1) the coupling is **Ising-ZZ** and *each qubit* belongs to a *type* that is jointly operator controllable locally,
- (2) the coupling is **Heisenberg-XXX** and *a single qubit* is controllable locally,
- (3) the coupling is **Heisenberg-XX** and *one adjacent qubit pair* is fully operator controllable ($\mathfrak{su}(4)$).



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Design Rules

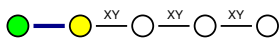
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Minimalistically Controlled Systems

Examples

quant-ph/0904.4654



$$H_d := \frac{1}{2} \sum_{k=2}^{n-1} (1 + \gamma) X_k X_{k+1} + (1 - \gamma) Y_k Y_{k+1} + \sum_{i=2}^n B_i Z_i$$

Example (inner symmetries)

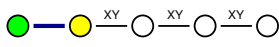
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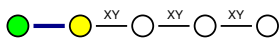
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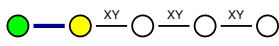
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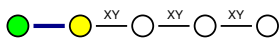
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Control of Markovian Open Systems

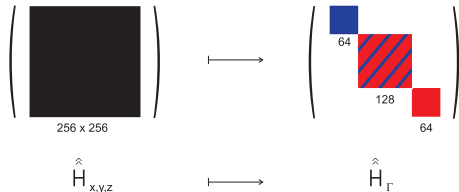
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Original Principle:

Code logical qubits in decoherence-free *physical* levels

- master equation: $\dot{\rho} = -(i \text{ad}_H + \Gamma) \rho$
- **DFS**: eigenspace to Γ with **eigenval = 0** (Bell states \mathcal{B})
- Express $\hat{H} \equiv \text{ad}_H$ in eigenbasis of Γ (here 4 qubits)



- Task: perform calculation (e.g. CNOT) **within DFS**

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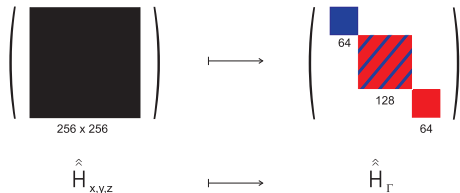
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Control of Markovian Open Systems

System of 2 Qubits Coded in 4 Spins

- 1 logical qubit coded by 2 physical qubits in Bell states

$$|0\rangle_L := |\psi^+\rangle = \frac{1}{\sqrt{2}}(|01\rangle + |10\rangle), \quad |1\rangle_L := |\psi^-\rangle = \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle)$$

$$\mathcal{B} := \text{span} \{ |\psi^\pm\rangle\langle\psi^\pm|, |\psi^\mp\rangle\langle\psi^\pm| \}$$

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Control of Markovian Open Systems

System of 2 Qubits Coded in 4 Spins

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- 2 logical qubits coded by 4 physical qubits

A diagram showing two qubits, represented by black dots, connected by a horizontal line. Above the line is the expression $\frac{1}{\sqrt{2}}(|01\rangle + |10\rangle)$ and below the line is $\frac{1}{\sqrt{2}}(|01\rangle - |10\rangle)$.

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Control of Markovian Open Systems

System of 2 Qubits Coded in 4 Spins

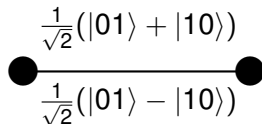
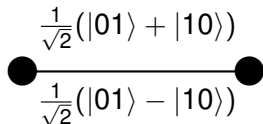
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- 2 logical qubits coded by 4 physical qubits



- protection against T_2 relaxation (Redfield: $\Gamma \sim [ZZ, [ZZ, \rho]]$)

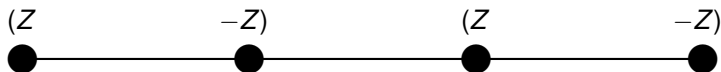
because $[\rho, ZZ] = 0 \quad \forall \quad \rho \in \mathcal{B} \otimes \mathcal{B}$



Control of Markovian Open Systems

Model with 4 Linearly Coupled Spins

■ controls



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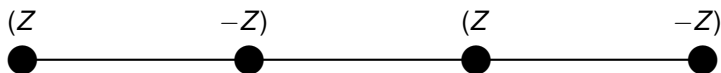
A3



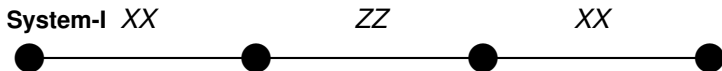
Control of Markovian Open Systems

Model with 4 Linearly Coupled Spins

■ controls



■ drift: Ising (ZZ) and Heisenberg (XX) interactions



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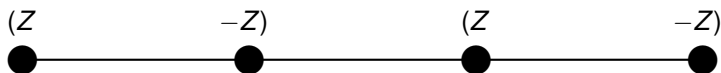
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Control of Markovian Open Systems

Model with 4 Linearly Coupled Spins

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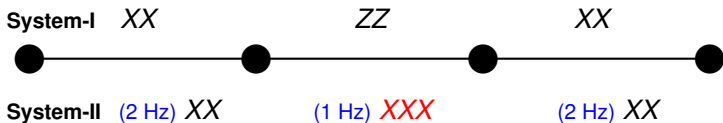
Control of Markovian Open Systems

Model with 4 Linearly Coupled Spins

■ controls



■ drift: Ising (ZZ) and Heisenberg (XX,XXX) interactions



■ relaxation ($T_2^{-1} : T_1^{-1} = 4.0 \text{ s}^{-1} : 0.024 \text{ s}^{-1} \simeq 170 : 1$)

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Control of Markovian Open Systems

Algebraic Analysis of System I

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■ System-I: staying **within** slowly-relaxing subspace

- drift Hamiltonian D_1 with **Ising-ZZ**
- controls $C_{1,2}$

$$D_1 := J_{xx} (xx\mathbb{1}\mathbb{1} + \mathbb{1}\mathbb{1}xx + yy\mathbb{1}\mathbb{1} + \mathbb{1}\mathbb{1}yy) + J_{zz} \mathbb{1}zz\mathbb{1}$$

$$C_1 := z\mathbb{1}\mathbb{1}\mathbb{1} - \mathbb{1}z\mathbb{1}\mathbb{1}$$

$$C_2 := \mathbb{1}\mathbb{1}z\mathbb{1} - \mathbb{1}\mathbb{1}\mathbb{1}z .$$

$$\Rightarrow \langle D_1, C_1, C_2 \rangle_{\text{Lie}} \Big|_{\mathcal{B} \otimes \mathcal{B}} \stackrel{\text{rep}}{=} \mathbf{su}(4)$$

- Liouville subspace $\mathcal{B} \otimes \mathcal{B}$ of Bell states
spans states protected against T_2 -relaxation



Control of Markovian Open Systems

Algebraic Analysis of System I

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Control of Markovian Open Systems

Algebraic Analysis of System II

■ System-II: driving **outside** slowly-relaxing subspace

- drift: extended to **isotropic Heisenberg-XXX**

$$D_1 + D_2 := J_{xx} (xx\mathbf{1}\mathbf{1} + \mathbf{1}\mathbf{1}xx + yy\mathbf{1}\mathbf{1} + \mathbf{1}\mathbf{1}yy) \\ + J_{xyz} (\mathbf{1}xx\mathbf{1} + \mathbf{1}yy\mathbf{1} + \mathbf{1}zz\mathbf{1})$$

- Lie-algebraic closure: in **66-dim. Lie algebra**

$$\dim\langle (D_1 + D_2), C_1, C_2 \rangle_{\text{Lie}} = 66 \quad \left[\stackrel{\text{iso}}{=} \mathfrak{so}(12) \right]$$

- **su(4)** merely **subalgebra**

$$\dim\langle D_1, C_1, C_2 \rangle_{\text{Lie}} = 15$$

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Control of Markovian Open Systems

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Control of Markovian Open Systems

Algebraic Analysis of System II

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Control of Markovian Open Systems

Algebraic Analysis of System II

System-II:

- full controllability **within** slowly-relaxing subspace

- observation

$$e^{-i\pi C_1} (D_1 + D_2) e^{i\pi C_1} = D_1 - D_2$$

- Trotter limit

$$\lim_{n \rightarrow \infty} \left(e^{-i(D_1 + D_2)/(2n)} e^{-i(D_1 - D_2)/(2n)} \right)^n = e^{-iD_1}$$

- reduction of dynamics

System-II $\xrightarrow{\text{infinite \# switchings}}$ System-I

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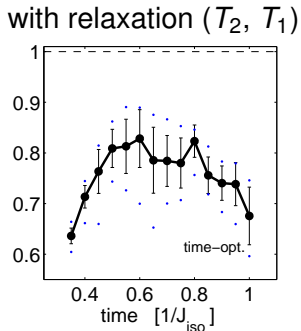
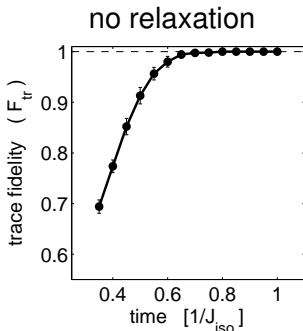
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Examples of Quantum Control

Decoherence Control: Results

Typical: system drives **outside** protected subspace



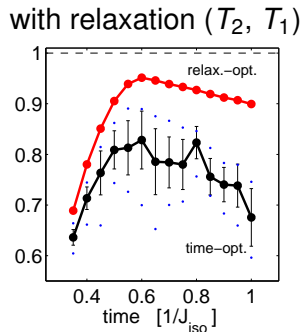
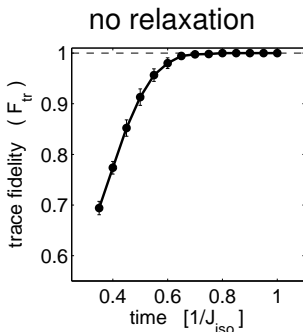
- mean of 15 time-optimised pulse sequences
- dissipation affects sequences differently



Examples of Quantum Control

Decoherence Control: Results

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- mean of 15 time-optimised pulse sequences
- dissipation affects sequences differently
- relaxation-optimised: **systematic substantial gain**

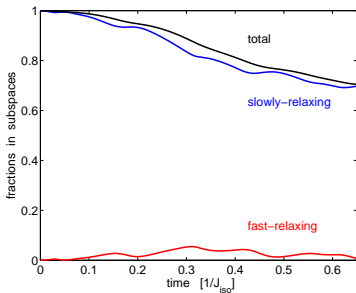


Control of Markovian Open Systems

Realising Quantum Gates with Minimal Relaxation

CNOT: Projection into Subspaces

■ time-optimised



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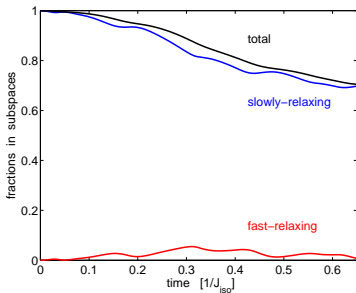


Control of Markovian Open Systems

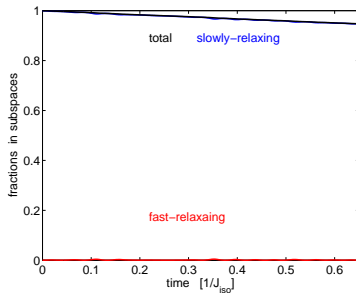
Realising Quantum Gates with Minimal Relaxation

CNOT: Projection into Subspaces

■ time-optimised



■ opt. against decoherence



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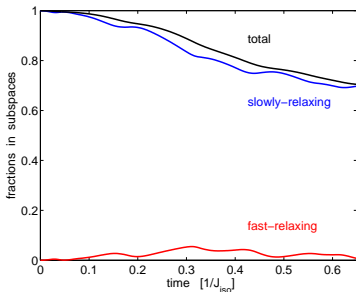


Control of Markovian Open Systems

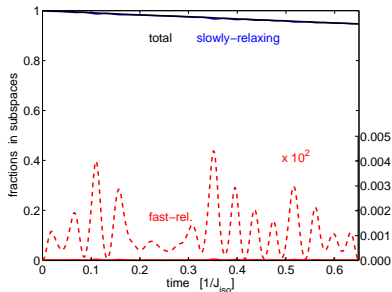
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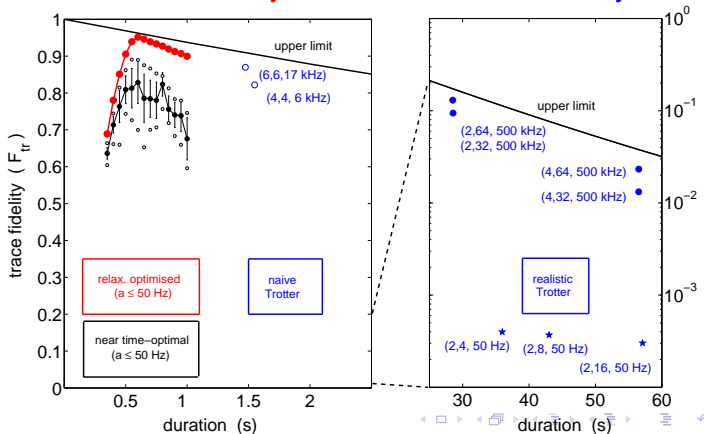
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Realising Quantum Gates with Minimal Relaxation

quant-ph/0609037

■ CNOT: comparison of methods

mbox **by decoherence control:** **> 95% fidelity** **conventional:** **< 15% fidelity**



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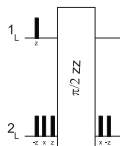


Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ logical qubits



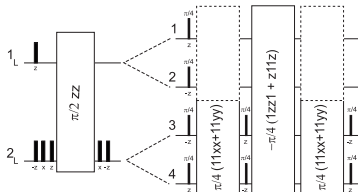
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Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ physical qubits



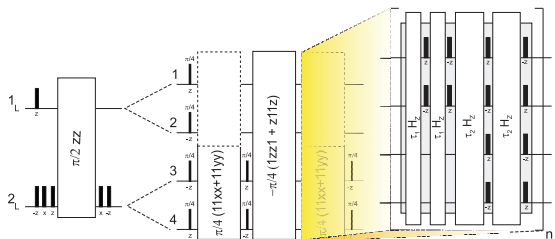
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Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ realisation by **System-I**



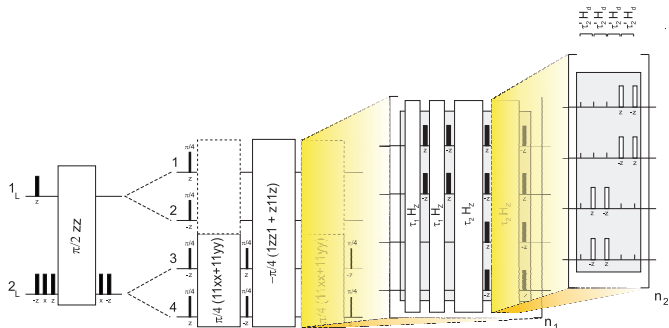
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Alternative Decoherence Control

Paper and Pen Approach: TROTTER Expansion

Decoherence-Protected CNOT-Gate via

■ realisation by **System-II**



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Control of Non-Markovian Open Systems

Qubit Coupled via Two-Level Fluctuator to Spin Bath

with P. Rebentrost and F. Wilhelm

Model:

qubit coupled to a two-level fluctuator coupled to a bath

$$H = H_S + H_I + H_B$$

- $H_S = E_1(t)\sigma_Z + \Delta\sigma_X + E_2\tau_Z + \Lambda\sigma_Z\tau_Z$

- $H_I = \sum_i \lambda_i (\tau^+ b_i + \tau^- b_i^\dagger)$

- $H_B = \sum_i \hbar\omega_i b_i^\dagger b_i$

Ohmic bath spectrum: $J(\omega) = \sum_i \lambda_i^2 \delta(\omega - \omega_i) = \kappa\omega\Theta(\omega - \omega_c)$

couplings λ_i , damping κ , high-freq. cut-off ω_c

PRL 102 090401 (2009)

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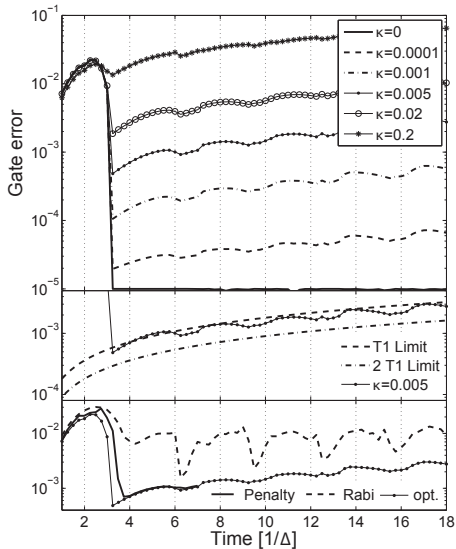


Control of Non-Markovian Open Systems

Qubit Coupled via Two-Level Fluctuator to Spin Bath

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← RABI pulse
← cut error by factor ≤ 10 with optimal control

Control of Non-Markovian Open Systems

- Principle: embed to **Markovian** and **project**

$$\begin{array}{ccc}
 \rho_0 = \rho_{SE}(0) \otimes \rho_B(0) & \xrightarrow{\text{Ad}_W(t)} & \rho(t) = W(t)\rho_0 W^\dagger(t) \\
 \Pi_{SE} \downarrow \text{tr}_B & & \Pi_{SE} \downarrow \text{tr}_B \\
 \rho_{SE}(0) & \xrightarrow[\text{Markovian}]{F_{SE}(t)} & \rho_{SE}(t) \\
 \Pi_S \downarrow \text{tr}_E & & \Pi_S \downarrow \text{tr}_E \\
 \rho_S(0) & \xrightarrow[\text{non-Markovian}]{F_S(t)} & \rho_S(t)
 \end{array}$$

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- Markovianity, Divisibility I
- Lie Semigroups
- GKS-Lindblad Gen.
- Divisibility II



■ Viewing Markovian Quantum Channels as Lie Semigroups with GKS-Lindblad Generators as Lie Wedge



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Observe: **two notions**

Definition

- A CP-Map T is *(infinitely) divisible*, if $\forall r \in \mathbb{N}$ there is a S with $T = S^r$.
- A CP-map T is *infinitesimally divisible* if $\forall \epsilon > 0$ there is a sequence $\prod_{j=1}^r S_j = T$ with $\|S_j - \text{id}\| \leq \epsilon$.



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Markovianity \Leftrightarrow Divisibility

Basic Structure

Notions:

time-(in)dependent CP-map: solution of
time-(in)dependent master eqn. $\dot{X} = -\mathcal{L} \circ X$.

Theorem (Wolf & Cirac (2008))

- *The set of all time-independent Markovian CP-maps coincides with the set of all (infinitely) divisible CP-maps.*
- *The set of all time-dependent Markovian CP-maps coincides with the closure of the set of all infinitesimally divisible CP-maps.*

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Markovianity \Leftrightarrow Divisibility

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Notions:

time-(in)dependent CP-map: solution of
time-(in)dependent master eqn. $\dot{X} = -\mathcal{L} \circ X$.

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Observe: **semigroup structure**

$$\text{Reach}(\mathbf{1}, t_1) \circ \text{Reach}(\mathbf{1}, t_2) = \text{Reach}(\mathbf{1}, t_1 + t_2) \quad \forall t_i \geq 0$$

Definition

- A **subsemigroup** $\mathbf{S} \subset \mathbf{G}$ of a Lie group \mathbf{G} with algebra \mathfrak{g} contains $\mathbf{1}$ and follows $\mathbf{S} \circ \mathbf{S} \subseteq \mathbf{S}$. Its largest subgroup is denoted $E(\mathbf{S}) := \mathbf{S} \cap \mathbf{S}^{-1}$.
- Its **tangent cone** is defined by

$$L(\mathbf{S}) := \{\dot{\gamma}(0) \mid \gamma(0) = \mathbf{1}, \gamma(t) \in \mathbf{S}, t \geq 0\} \subset \mathfrak{g},$$

for any $\gamma : [0, \infty) \rightarrow \mathbf{G}$ being a smooth curve in \mathbf{S} .

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Definition (Lie Wedge and Lie Semialgebra)

- A *wedge* \mathfrak{w} is a closed convex cone of a finite-dim. real vector space.
- Its *edge* $E(\mathfrak{w}) := \mathfrak{w} \cap -\mathfrak{w}$ is the largest subspace in \mathfrak{w} .
- It is a *Lie wedge* if it is invariant under conjugation

$$e^{\text{ad}_g}(\mathfrak{w}) \equiv e^g \mathfrak{w} e^{-g} = \mathfrak{w}$$

for all edge elements $g \in E(\mathfrak{w})$.

- A *Lie semialgebra* is a Lie wedge compatible with BCH multiplication $X * Y := X + Y + \frac{1}{2}[X, Y] + \dots$ so that for a BCH neighbourhood B of $0 \in \mathfrak{g}$

$$(\mathfrak{w} \cap B) * (\mathfrak{w} \cap B) \in \mathfrak{w} .$$



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Define as completely positive, trace-preserving invertible linear operators the set \mathbf{P}^{cp} , and let \mathbf{P}_0^{cp} denote the connected component of the unity.

Theorem (Kossakowski, Lindblad)

The Lie wedge to the connected component of the unity of the semigroup of all invertible CPTP maps is given by the set of all linear operators of GKS-Lindblad form:

$$L(\mathbf{P}_0^{\text{cp}}) = \{-\mathcal{L} \mid \mathcal{L} = -(i \text{ad}_H + \Gamma_L)\} \quad \text{with}$$

$$\Gamma_L(\rho) = \frac{1}{2} \sum_k \{V_k^\dagger V_k, \rho\}_+ - 2V_k \rho V_k^\dagger$$

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Theorem

The semigroup

$$\mathbf{T} := \overline{\langle \exp(\mathbf{L}(\mathbf{P}_0^{\text{cp}})) \rangle_S} \subseteq \mathbf{P}_0^{\text{cp}}$$

generated by $\mathbf{L}(\mathbf{P}_0^{\text{cp}})$ is a Lie subsemigroup with global Lie wedge $\mathbf{L}(\mathbf{T}) = \mathbf{L}(\mathbf{P}_0^{\text{cp}})$.



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Corollary (to Wolf, Cirac (2008))

\mathbf{P}_0^{cp} itself is *not a Lie subsemigroup*, yet it comprises

- (1) the set of *time independent Markovian channels*, i.e. the union of all one-parameter Lie semigroups $\{\exp(-\mathcal{L}t) \mid t \geq 0\}$ with \mathcal{L} in GKS-Lindblad form;
- (2) the *closure* of the set of *time dependent Markovian channels*, i.e. the Lie semigroup \mathbf{T} ;
- (3) a set of *non-Markovian channels* whose intersection with \mathbf{P}_0^{cp} has non-empty interior.



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Corollary

A quantum channel is *time dependent Markovian* iff it allows for a representation $T = \prod_{j=1}^r S_j$, where $S_1 = e^{-\mathcal{L}_1}, S_2 = e^{-\mathcal{L}_2}, \dots, S_r = e^{-\mathcal{L}_r}$ so that there is a global *Lie wedge* \mathfrak{w}_r generated by $\mathcal{L}_1, \mathcal{L}_2, \dots, \mathcal{L}_r$.



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Corollary

Let $T = \prod_{j=1}^r S_j$ be a *time dependent Markovian channel* with $S_1 = e^{-\mathcal{L}_1}$, $S_2 = e^{-\mathcal{L}_2}$, \dots , $S_r = e^{-\mathcal{L}_r}$ and let \mathfrak{w}_r denote the smallest global Lie wedge generated by $\mathcal{L}_1, \mathcal{L}_2, \dots, \mathcal{L}_r$. Then

- T boils down to a *time independent Markovian channel*, if it is sufficiently close to the unity and if there is a representation so that the associated Lie wedge \mathfrak{w}_r specialises to a *Lie semialgebra*.

Complements recent work: Wolf, Cirac, *Commun. Math. Phys.* (2008) & Wolf, Eisert, Cubitt, Cirac, *PRL* (2008)



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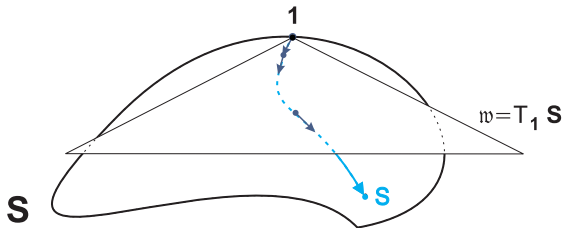
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Consider: controlled system with *time dep* Liouvillians $\{\mathcal{L}_u(t)\}$

$$\dot{X} = -\mathcal{L}_u(t)X = -(iH_d + i\sum_j u_j(t)H_j + \Gamma)X$$



Liouvillians \mathcal{L}_u form

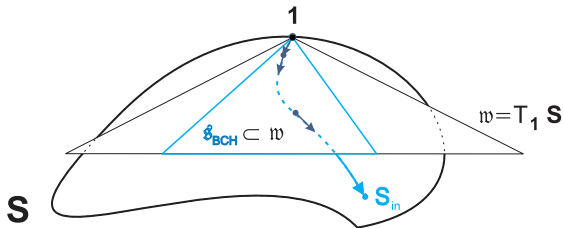
- *Lie wedge* \mathfrak{w}
- *Lie semialgebra* $\mathfrak{s} \subset \mathfrak{w}$ if $\{\mathcal{L}_u\}$ BCH compatible with \mathfrak{w}
 then $\{e^{-t\mathcal{L}_{\text{eff}}} \mid t > 0\}$ **physical** at all times.
- Else $\{e^{-t\mathcal{L}_{\text{eff}}} \mid t > 0\}$ **unphysical** except $t = 0$; $t = t_{\text{eff}}$ etc.

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Liouvillians \mathcal{L}_u form

$$L_j * L_k := L_j + L_k + \frac{1}{2}[L_j, L_k] + \dots \in w$$

■ *Lie wedge* w

■ *Lie semialgebra* $s \subset w$ if $\{\mathcal{L}_u\}$ BCH compatible with w

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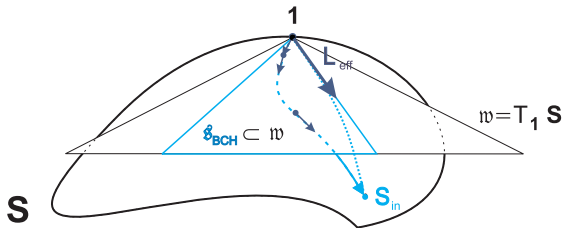
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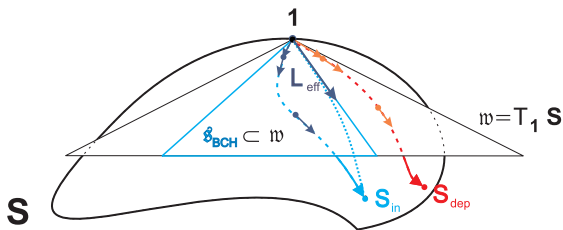
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