

Evaluation of convex roof entanglement measures

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1 Motivation

- Why entanglement quantification is important?

2 Calculating entanglement measures

- Convex roof of the entropy
- Tangle
- Other quantities
- Even tighter lower bounds

Why entanglement quantification is important?

- Many experiments are aiming to create entangled states.
- We need to calculate entanglement measures for these states.
- Apart from trivial system sizes, we cannot do it.

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Entanglement

The entanglement of a bipartite quantum state

For pure states, it is defined as

$$E(|\Psi\rangle) = S[\text{Tr}_1(|\Psi\rangle)],$$

for pure states, where S is an entropy.

For mixed states, it is defined with a **convex roof** as

$$E(\rho) = \min_{\{p_k, |\Psi_k\rangle\}} \left(\sum_k p_k E(|\Psi_k\rangle) \right),$$

where $\{p_k, |\Psi_k\rangle\}$ is a decomposition to pure states

$$\rho = \sum_k p_k |\Psi_k\rangle\langle\Psi_k|.$$

Linear entropy for pure states

- Linear entropy

$$S_{\text{lin}}(\rho) = 1 - \text{Tr}(\rho^2).$$

- Known: linear entropy of entanglement for pure states can be defined as an expectation value on **two copies** (AB and $A'B'$) as

$$E_{\text{lin}}(|\Psi\rangle) = \text{Tr}[\mathcal{A}_{AA'} \otimes \mathbb{1}_{BB'} (|\Psi\rangle\langle\Psi|)_{AB} \otimes (|\Psi\rangle\langle\Psi|)_{A'B'}],$$

where

$$\mathcal{A}_{AA'} := (\mathbb{1} - \mathcal{F})_{AA'}$$

and \mathcal{F} is the flip operator.

Linear entropy for mixed states: convex roof

- For mixed states

$$\begin{aligned} E_{\text{lin}}(\rho) &= \min_{\{\rho_k, |\psi_k\rangle\}} \sum_k \rho_k E_{\text{lin}}(|\psi_k\rangle) = \\ &= \min_{\{\rho_k, |\psi_k\rangle\}} \sum_k \rho_k \text{Tr}(\mathcal{A}_{AA'} |\psi_k\rangle\langle\psi_k|^{\otimes 2}) \\ &= \min_{\omega_{12}} \text{Tr}(\mathcal{A}_{AA'} \omega_{12}), \end{aligned}$$

where ω_{12} are symmetric separable states, i.e.,

$$\omega_{12} = \sum_k \rho_k |\psi_k\rangle\langle\psi_k| \otimes |\psi_k\rangle\langle\psi_k|.$$

- This is the key step in our approach.

Surprise 1

- Mapping of the problem

Optimization over decompositions \longrightarrow Optimization over symmetric separable states

- We connected the separability theory to a general mathematical problem.

How to calculate it

- The convex roof of the linear entropy can be written as

$$\begin{aligned} E_{\text{lin}}(\varrho) = \min_{\omega_{12}} & \quad \text{Tr}(\mathcal{A}_{AA'}\omega_{12}), \\ \text{s.t.} & \quad \omega_{12} \text{ is symmetric, separable,} \\ & \quad \omega_1 = \varrho, \end{aligned}$$

where $\omega_1 \equiv \text{Tr}_2(\omega_{12})$.

- A lower bound can be obtained as with the PPT condition

$$\begin{aligned} E_{\text{lin}}(\varrho) = \min_{\omega_{12}} & \quad \text{Tr}(\mathcal{A}_{AA'}\omega_{12}), \\ \text{s.t.} & \quad \omega_{12} \text{ is symmetric PPT,} \\ & \quad \omega_1 = \varrho, \end{aligned}$$

where $\omega_1 \equiv \text{Tr}_2(\omega_{12})$. **This is a semidefinite program.**

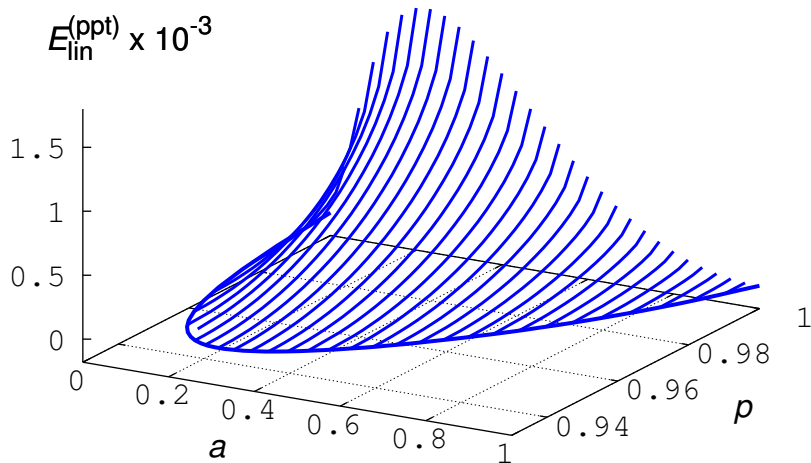
Surprise 2

- The lower bound
 - is nonzero for all states with a non-positive semidefinite partial transpose (NPPT).
 - is nonzero for some states with a positive semidefinite partial transpose (PPT).
- For all non-PPT states and for all states that do not have a $2 : 2$ symmetric extension we have a nonzero bound.
- Moreover, for all states having a $2:2$ PPT symmetric extension the bound is zero.

[Extensions: Doherty, Parrilo, Spedalieri, PRA 69, 022308 (2004)]

Example: Entanglement of a PPT state

- 3×3 Horodecki state mixed with white noise.
- a = parameter of the state, $1 - p$ = noise fraction



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Wootters' Tangle

- The well-known tangle for three-qubits can be defined as a fourth-order polynomial in expectation values.
[A. Osterloh and J. Siewert, Phys. Rev. A 86, 042302 (2012).]
- Hence, it can be obtained as an optimization over four-partite symmetric separable states

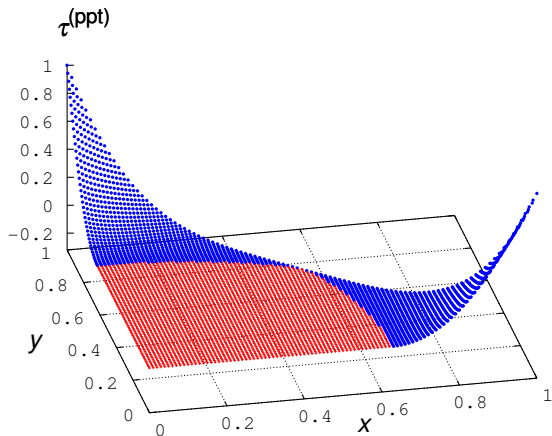
$$\begin{aligned} \tau(\varrho) = \min_{\omega_{1234}} & \quad \text{Tr}(T\omega_{1234}), \\ \text{s.t.} & \quad \omega_{1234} \text{ symmetric, fully separable,} \\ & \quad \omega_1 = \varrho, \end{aligned}$$

where T is an operator (4 parties with 3 qubits each).

- Similar idea works: replace separable states by PPT states.

Example: tangle of a two-parameter family of states

$$\rho(x, y) = x|\text{GHZ}^+\rangle\langle\text{GHZ}^+| + y|\text{GHZ}^-\rangle\langle\text{GHZ}^-| + (1 - x - y)|\text{W}\rangle\langle\text{W}|$$



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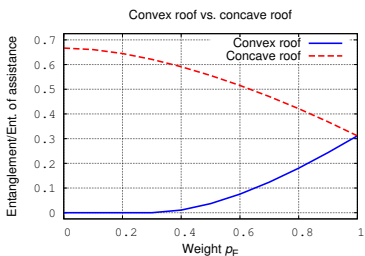
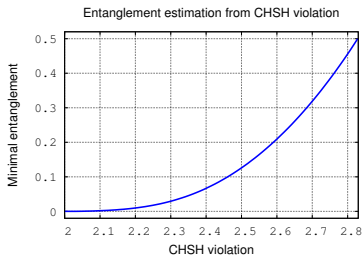
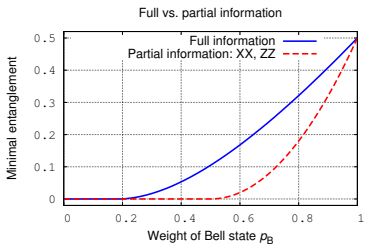
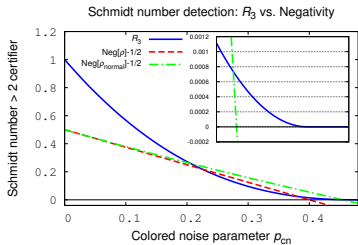
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Other quantities

- **Schmidt number**. I.e., the convex roof of $R_3(|\Psi\rangle) = \sum_{i<j<k} \lambda_i \lambda_j \lambda_k$ tells us whether the Schmidt number is larger than 2.
- Entanglement vs. **CHSH violation**
- Lower bound on entanglement **based on some measurement results**
- **Concave roof** instead of convex roofs: E. of assistance
- Lower bound on **quantum Fisher information** based some measurement results.
[Tóth, Petz, PRA 2013.]
- One can get even a **witness**!!

[For references, please see our work on the arXiv.]

Examples



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A series of tighter and tighter lower bounds

- To strengthen the bound, a criterion stronger than PPT must be employed.
- For example, the method of PPT symmetric extensions can be used.

[Doherty, Parrilo, Spedalieri, *Phys. Rev. A* 69, 022308 (2004)]

- Sequence of lower bounds $E_{\text{lin}}^{(n)}$ with increasing accuracies.
- Calculation: semidefinite program.

Summary

- We showed how to obtain a good lower bound on quantities defined with convex roofs.
- We used it for calculating entanglement measures and the tangle, and several other quantities.

See:

GT, T. Moroder, and O. Gühne,
Evaluation of convex roof entanglement measures,
Phys. Rev. Lett, in press; arxiv:1409.3806.

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THANK YOU FOR YOUR ATTENTION!

