## General characteristics of multi-partite quantum systems

 (Lecture of the Quantum Information class of the Master in Quantum Science and Technology)Géza Tóth

Theoretical Physics, University of the Basque Country (UPV/EHU), Bilbao, Spain
Donostia International Physics Center (DIPC), San Sebastián, Spain
IKERBASQUE, Basque Foundation for Science, Bilbao, Spain Wigner Research Centre for Physics, Budapest, Hungary

UPV/EHU, Leioa
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## Outline

(1) General characteristics of multi-partite quantum systems

- A. Classical bits
- B. Quantum bit - pure states
- C. Multi-qubit systems - pure states
- D. Measurement
- E. Mixed states and the density matrix
- F. Geometry of quantum states
- A single qubit
- A single qudit (qunit):d-dimensional systems
- G. Two or more qubits: reduced states
- H. Purifications
- I. Purity
- J. Entropy
- Shannon entropy
- Von Neumann entropy
- Quantum relative entropy
- Linear entropy
- K. Fidelity


## A single classical bit

- A classical bit can be either 0 or 1 . Can we still use it to describe a real number between 0 and 1?
- For that, we need an ensemble of several classical bits

$$
\begin{equation*}
\left\{b_{k}\right\}_{k=1}^{M} \tag{1}
\end{equation*}
$$

where $b_{k}=0$ or 1

- We can interpret the average value and the variance. That is,

$$
\begin{equation*}
\langle b\rangle=\frac{1}{M} \sum_{k} b_{k} \tag{2}
\end{equation*}
$$

and

$$
\begin{equation*}
(\Delta b)^{2}=\frac{1}{M} \sum_{k}\left(b_{k}-\langle b\rangle\right)^{2} \tag{3}
\end{equation*}
$$

## A single classical bit II

- This can also be given with probabilities:
- Let $P_{0}$ and $P_{1}$ be the probabilities of having a 0 or a 1 .
- The expectation value and the variance are the function of $P_{0}$ and $P_{1}$. Since $P_{0}+P_{1}=1$, we have a single real degree of freedom that describes the statistical properties of an ensemble of bits.
- Hence,

$$
\begin{equation*}
\langle b\rangle=P_{1} \tag{4}
\end{equation*}
$$

and

$$
\begin{equation*}
(\Delta b)^{2}=P_{0}\left(0-P_{1}\right)^{2}+P_{1}\left(1-P_{1}\right)^{2} \tag{5}
\end{equation*}
$$

## Stochastic computing

- Stochastic computing uses random bits to calculate (John von Neumann, 1953).
- A random bit represents a real number between 0 and 1. Two random bits can easily be multiplied.

$$
\begin{equation*}
\left\langle b_{1} b_{2}\right\rangle=\left\langle b_{1}\right\rangle\left\langle b_{2}\right\rangle \tag{6}
\end{equation*}
$$

- We need many samples to get the average with small error.


## Stochastic computing II

Lectures on
PROBABILISTIC LOGICS AND THE SYNTHESIS OF RELIABLE ORGANISMS FROM UNRELIABLE COMPONENTS

delivered by

PROFESSOR J. von NEUMANN
The Institute for Advanced Study
Princeton, N. J.

## Stochastic computing III



The RASCEL stochastic computer, circa 1969, Wikipedia.

## Stochastic computing IV

## Multiplication is possible with an AND gate.



$$
\begin{aligned}
& 11001110000010010011 \\
& 11101110000000000011
\end{aligned}
$$

Figure 1.2: Similarity of biological signals and stochastic numbers; information is carried via pulses.


Figure 1.3: Stochastic multiplication: (a) accurate result with uncorrelated inputs; (b) inaccurate result due to correlated inputs.
A. Alaghi, The Logic of Random Pulses: Stochastic Computing,

Ph.D. Thesis, University of Michigan, 2015.

## Several classical bits

- $N$ classical bits can be in one of the $2^{N}$ binary states. For example, for $N=2$, these are $00,01,10$ and 11 .
- For $N=2$, these are

$$
\begin{equation*}
P_{00}, P_{01}, P_{10}, P_{11} \tag{7}
\end{equation*}
$$

- The ensemble of the $N$-bit units can be described by the $2^{N}$ probabilities.
- Since, again, the sum of all the probablities is 1 , we need $2^{N}-1$ real degrees of freedom to describe the statistical properties of such an ensemble.


## Several classical bits II

- Let us consider some function of $N$ bits $f(k)$, where $k$ is now an $N$ bit number.
- Then, the expectation value of $f$ is

$$
\begin{equation*}
\langle f\rangle=\sum_{k=0}^{2^{N}-1} p_{k} f(k)=\vec{p} \vec{f} \tag{8}
\end{equation*}
$$

where $k$ is an $N$-bit number, i.e., an integer between 0 and $2^{N}-1$. We put the $f_{k}$ 's into a vector $\vec{f}$. We also put the $p_{k}$ probabilities into $\vec{p}$.

## Several classical bits III

- We can also write

$$
\begin{equation*}
\left\langle f^{2}\right\rangle=\sum_{k} p_{k} f_{k}^{2} \tag{9}
\end{equation*}
$$

Hence,

$$
\begin{equation*}
(\Delta f)^{2}=\sum_{k} p_{k} f_{k}^{2}-\left(\sum_{k} p_{k} f_{k}\right)^{2} \tag{10}
\end{equation*}
$$

These were relevant, since in the quantum case, we will have similar expressions.

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## Quantum bit - pure states

- A quantum bit (=two-state system, spin- $\frac{1}{2}$ particle) can be in a pure state

$$
\begin{equation*}
|q\rangle=\alpha|0\rangle+\beta|1\rangle, \tag{11}
\end{equation*}
$$

where $\alpha$ and $\beta$ are complex numbers, and the normalisation condition $|\alpha|^{2}+|\beta|^{2}=1$.

- Note that the overall phase does not matter, thus a pure quantum bit is described by two degrees of freedom.
- The two complex coefficients have 4 real degrees of freedom.
- However, due to the normalisation condition and the arbitrariness of the overall phase we are left with two degrees of freedom.)


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## Multi-qubit systems - pure states

- What about a two-qubit system? What kind of states it can be in? One could think on qubit 1 in state

$$
\begin{equation*}
\left|q_{1}\right\rangle=\alpha_{1}|0\rangle+\beta_{1}|1\rangle \tag{12}
\end{equation*}
$$

and qubit 2 in state

$$
\begin{equation*}
\left|q_{2}\right\rangle=\alpha_{2}|0\rangle+\beta_{2}|1\rangle . \tag{13}
\end{equation*}
$$

- However, we all know that the general state of the two-qubit system can be given as

$$
\begin{equation*}
\left|q_{12}\right\rangle=\alpha_{00}|00\rangle+\alpha_{01}|01\rangle+\alpha_{10}|00\rangle+\alpha_{11}|01\rangle . \tag{14}
\end{equation*}
$$

## Multi-qubit systems - pure states II

- In general, for $N$ qubits we need $N$ complex numbers. Again the state has to be normalized and the overall phase does not matter, thus this means $2 \times 2^{N}-2$ real degrees of freedom.
- We can place the coefficients in a vector, called state vector and write

$$
|\Psi\rangle=\left(\begin{array}{l}
\alpha_{00}  \tag{15}\\
\alpha_{01} \\
\alpha_{10} \\
\alpha_{11}
\end{array}\right) .
$$

- The properties of the state vector are: it is normalized

$$
\begin{equation*}
\langle\Psi \mid \Psi\rangle=1 . \tag{16}
\end{equation*}
$$

## Multi-qubit systems - pure states III

- An overall phase does not matter:

$$
\begin{equation*}
e^{-i \theta}|\Psi\rangle \tag{17}
\end{equation*}
$$

describes the same state for any $\theta$.

- The expectation value of an operator for a pure state can be obtained as

$$
\begin{equation*}
\langle A\rangle=\langle\Psi| A|\Psi\rangle=\operatorname{Tr}(A|\Psi\rangle\langle\Psi|) . \tag{18}
\end{equation*}
$$

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## Measurement

- The von Neumann measuement in the $z$ basis results is eithet 0 or 1. If the state was $\alpha|0\rangle+\beta|1\rangle$, then we get a statistical mixture of 0 and 1 , with the probabilities

$$
\begin{equation*}
P_{0}=|\alpha|^{2} \tag{19}
\end{equation*}
$$

and

$$
\begin{equation*}
P_{1}=|\beta|^{2} \tag{20}
\end{equation*}
$$

That is, from an ensemble of quantum bits we get an ensemble of classical bits.

- If we measure in the $x$ basis, we get another classical ensemble.
- For a multi-qubit system, if we measure in the some basis (e.g., $x, y$ or $z$ ), we get an ensemble of $N$-bit systems. However, for exach choice of basis we get a different classical ensemble.


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## Mixed states and the density matrix

- So far we were talking about pure states.
- In reality, in an experiment we do not have a situation where a machine always produces the $\left|\Psi_{1}\right\rangle$ state.
- Sometimes it makes mistakes, and produces the $\left|\Psi_{k}\right\rangle$ states for $k=2,3, \ldots$ How to describe such a situation?

| $\left\|\Psi_{1}\right\rangle$ | $p_{1}$ |
| :---: | :---: |
| $\left\|\Psi_{2}\right\rangle$ | $p_{2}$ |
| $\left\|\Psi_{3}\right\rangle$ | $p_{3}$ |
| $\ldots$ | $\ldots$ |

## Mixed states and the density matrix

- What is the expectation value of an operator in such a system?

We can write it as

$$
\begin{equation*}
\langle A\rangle=\sum_{k} p_{k}\left\langle\Psi_{k}\right| A\left|\Psi_{k}\right\rangle=\operatorname{Tr}\left(A \sum_{k} p_{k}\left|\Psi_{k}\right\rangle\left\langle\Psi_{k}\right|\right) \tag{21}
\end{equation*}
$$

- This can be rewritten as

$$
\begin{equation*}
\langle A\rangle=\operatorname{Tr}(\varrho A) \tag{22}
\end{equation*}
$$

where

$$
\begin{equation*}
\varrho=\sum_{k} p_{k}\left|\Psi_{k}\right\rangle\left\langle\Psi_{k}\right| \tag{23}
\end{equation*}
$$

is the density matrix (Neumann, Landau).

- Note that if $\varrho$ is diagonal, we obtain

$$
\begin{equation*}
\langle A\rangle=\operatorname{Tr}(\varrho A)=\sum_{k} \varrho_{k k} A_{k k} . \tag{24}
\end{equation*}
$$

That is, $A$ is written in the eigenbasis of $\varrho$. This is the scalar product of two vectors as in $\langle f\rangle=\vec{p} \vec{f}$ [given in Eq. (8)].

## Mixed states and the density matrix II

- The density matrix describes the state completely. Now we see, why the overall phase does not matter:

$$
\begin{equation*}
e^{-i \theta}\left|\Psi_{k}\right\rangle\left\langle\Psi_{k}\right| e^{+i \theta}=\left|\Psi_{k}\right\rangle\left\langle\Psi_{k}\right| \tag{25}
\end{equation*}
$$

- The properties of the density matrix are

$$
\begin{align*}
\varrho & =\varrho^{\dagger} \\
\varrho & \geq 0 \\
\operatorname{Tr}(\varrho) & =1 \tag{26}
\end{align*}
$$

- A $2^{N} \times 2^{N}$ density matrix has $4^{N}-1$ real parameters.
- For $N=1$, this means 3 real parameters, corresponding to the three coordinates of the Bloch vector. For $r N=2$, this means 8 real parameters.


## Mixed states and the density matrix III

- We can also say that

$$
\begin{equation*}
\operatorname{Tr}\left(\varrho^{2}\right) \leq 1 \tag{27}
\end{equation*}
$$

It is one only for pure (rank-1) states.

- The density matrix can be decomposed into the sum of pure states in many ways. The decomposition

$$
\begin{equation*}
\varrho=\sum_{k} p_{k}\left|\Psi_{k}\right\rangle\left\langle\Psi_{k}\right| \tag{28}
\end{equation*}
$$

is not unique, i.e., it is not necessarily an eigendecomposition. This has a large importance for entanglement theory.

## Mixed states and the density matrix IV

Summary:

|  | $N$ bits | $N$ qubits |
| :--- | :--- | :--- |
| Number of DOF | $2^{N}-1$ | $4^{N}-1$ |
| Description | $\vec{p}$ | $\varrho$ |
| Expectation value | $\vec{f} \vec{p}$ | $\operatorname{Tr}(A \varrho)$ |
| Normalization | $\sum_{k} p_{k}=1$ | $\operatorname{Tr}(\varrho)=1$ |

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## Bloch vector

- For a single qubit, the density matrix has three real parameters. It can be written as

$$
\begin{equation*}
\varrho=\frac{1}{2}\left(\mathbb{1}+\sum_{l=x, y, z} v_{l} \sigma_{l}\right), \tag{29}
\end{equation*}
$$

where $\sigma_{l}$ are the Pauli spin matrices.

- Using $\operatorname{Tr}\left(\sigma_{k} \sigma_{l}\right)=2 \delta_{k l}$, we can write

$$
\begin{equation*}
\operatorname{Tr}\left(\varrho^{2}\right)=\frac{1}{2}+\frac{1}{2} \sum_{I=x, y, z} v_{l}^{2} \tag{30}
\end{equation*}
$$

That is, the Bloch vector has a maximal length for pure states.

## Bloch vector II

- From $\operatorname{Tr}\left(\varrho^{2}\right) \leq 1$, the condition for being physical is Eq. (26), which is equalent to

$$
\begin{equation*}
\sum_{l=x, y, z}\left|v_{l}\right|^{2} \leq 1 \tag{31}
\end{equation*}
$$

The three-element vector is called the Bloch vector.

## Bloch vector III

- Let us identify the points in ( $v_{x}, v_{y}, v_{z}$ ) corresponding to physical states. They are in a ball.
- The pure states are on the surface.
- Mixed states are inside the Ball. This is because $\operatorname{Tr}\left(\varrho^{2}\right)$ is directly related to the length of the Bloch vector.
- The $|0\rangle$ and $|1\rangle$ correspond to the North and South Pole.
- $|0\rangle+\exp (-i \phi)|1\rangle$ correspond to points on the equator.


Set of physical quantum states for a single qubit. The axes correspond to $v_{l}$ for $I=x, y, z$. Pure states correspond to points on the surface, mixed states correspond to internal points.

## A single qudit (qunit):d-dimensional systems

- For higher dimensional systems the picture is much more complicated. Let us consider qudits with dimension $d$.
- Similarly to the case before, a $d \times d$ Hermitian matrix with a unit trace has $d^{2}-1$ degrees of freedom.
- Hence, we can write a density matrix as a linear combination of $d^{2}-1 \mathrm{SU}(\mathrm{d})$ generators as

$$
\begin{equation*}
\varrho=\frac{1}{d} \mathbb{1}+\frac{1}{2} \sum_{l=1}^{d^{2}-1} v_{l} g_{l} \tag{32}
\end{equation*}
$$

Here,

$$
\begin{equation*}
\operatorname{Tr}\left(g_{k} g_{l}\right)=2 \delta_{k l} . \tag{33}
\end{equation*}
$$

(Like for the Pauli matrices. Thus, we have something like the generalized Pauli matrices. $d=3$ : Gell-Mann matrices.)

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Here,

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\begin{equation*}
\operatorname{Tr}\left(g_{k} g_{l}\right)=2 \delta_{k l} \tag{35}
\end{equation*}
$$

- Like for the Pauli matrices. Thus, we have something like the generalized Pauli matrices. $d=3$ : for instance, Gell-Mann matrices.


## A single qudit (qunit):d-dimensional systems II

- Gell-Mann matrices:

$$
\begin{array}{ll}
\lambda_{1}=\left(\begin{array}{lll}
0 & 1 & 0 \\
1 & 0 & 0 \\
0 & 0 & 0
\end{array}\right) & \lambda_{2}=\left(\begin{array}{ccc}
0 & -i & 0 \\
i & 0 & 0 \\
0 & 0 & 0
\end{array}\right) \quad \lambda_{3}=\left(\begin{array}{ccc}
1 & 0 & 0 \\
0 & -1 & 0 \\
0 & 0 & 0
\end{array}\right) \\
\lambda_{4}=\left(\begin{array}{lll}
0 & 0 & 1 \\
0 & 0 & 0 \\
1 & 0 & 0
\end{array}\right) & \lambda_{5}=\left(\begin{array}{ccc}
0 & 0 & -i \\
0 & 0 & 0 \\
i & 0 & 0
\end{array}\right) \\
\lambda_{6}=\left(\begin{array}{lll}
0 & 0 & 0 \\
0 & 0 & 1 \\
0 & 1 & 0
\end{array}\right) & \lambda_{7}=\left(\begin{array}{ccc}
0 & 0 & 0 \\
0 & 0 & -i \\
0 & i & 0
\end{array}\right) \quad \lambda_{8}=\frac{1}{\sqrt{3}}\left(\begin{array}{ccc}
1 & 0 & 0 \\
0 & 1 & 0 \\
0 & 0 & -2
\end{array}\right) .
\end{array}
$$

There are other possibilities: J. Lawrence, quant-ph/0403095.

## A single qudit (qunit):d-dimensional systems III

- Let us again look at the points $\left(v_{1}, v_{2}, \ldots, v_{d^{2}-1}\right)$ corresponding to physical states.
- First note that the set of convex. This is because mixing two physical states $\varrho_{1}$ and $\varrho_{2}$, we always get a physical state

$$
\begin{equation*}
\varrho=p \varrho_{1}+(1-p) \varrho_{2} . \tag{36}
\end{equation*}
$$

## A single qudit (qunit):d-dimensional systems IV



Two convex objects and one that is not convex.

## A single qudit (qunit): $d$-dimensional systems V

- On the next figure we will show the set of quantum states.
- The cooridnate axis could be the $v_{l}$, for example.
- Inside the set there are the density matrices with full rank.
- On the boundary there are the states with less than full rank, such as for example rank-1 states, which are pure states.


Set of physical quantum states. Note that the set is convex. A,B,D: rank-1 states. C: rank-2 state. E: full rank states.

## A single qudit (qunit):d-dimensional systems VI

- Observation. The following inequality is true

$$
\begin{equation*}
\lambda_{\min }(A+B) \geq \lambda_{\min }(A)+\lambda_{\min }(B) \tag{37}
\end{equation*}
$$

Proof. Let us consider that for a Hermitian matrix $X$ we have

$$
\begin{equation*}
\lambda_{\min }(X)=\min _{\psi}\langle\psi| X|\psi\rangle . \tag{38}
\end{equation*}
$$

Then, for $A$ and $B$ Hermitian matrices we have

$$
\begin{align*}
\lambda_{\min }(A+B) & =\min _{\psi}\langle\psi| A+B|\psi\rangle \geq \min _{\psi}\langle\psi| A|\psi\rangle+\min _{\psi}\langle\psi| B|\psi\rangle \\
& =\lambda_{\min }(A)+\lambda_{\min }(B) . \tag{39}
\end{align*}
$$

We can prove similarly that

$$
\begin{equation*}
\lambda_{\max }(A+B) \leq \lambda_{\max }(A)+\lambda_{\max }(B) \tag{40}
\end{equation*}
$$

## Full rank states

- Using this, we can say the following.
- Observation. Full-rank states are inside the set. Proof. If the state is full rank, it means that for some small $\epsilon$

$$
\begin{equation*}
\varrho^{\prime}=\varrho+\epsilon H \tag{41}
\end{equation*}
$$

is also physical, where $H$ is a trace 0 Hermitian matrix. Why is that? See also the next figure.

## Full rank states II

## $\rho$ <br> - $\rho^{\prime}$

We take an internal state $\varrho$ and consider the states $\varrho^{\prime}$ in its neighborhood.

## Full rank states III

- It is physical since
(1) Trace is 1 .
(2) Hermitian.
(3) Eigenvalues are nonzero for small epsilon. This is because

$$
\begin{equation*}
\lambda_{\max }(\varrho)+\lambda_{\max }(\epsilon H) \geq \lambda_{k}\left(\varrho^{\prime}\right) \geq \lambda_{\min }(\varrho)+\lambda_{\min }(\epsilon H) . \tag{42}
\end{equation*}
$$

Here we have

$$
\lambda_{\min }(\epsilon H)= \begin{cases}+\epsilon \lambda_{\min }(H), & \text { if } \epsilon \geq 0,  \tag{43}\\ -|\epsilon| \lambda_{\max }(H), & \text { if } \epsilon<0 .\end{cases}
$$

Similar statement holds for $\lambda_{\max }(\epsilon H)$.

## Non-full-rank states

- Observation. Non-full-rank states are on the surface of the set.
- Proof. If the state is not full rank, then it has zero eigenvalues. Thus, there is an $H$ such that $\varrho^{\prime}$ is aphisical for any $\epsilon>0$ or any $\epsilon<0$.
- To be more explicit, let us write

$$
\begin{equation*}
\varrho=U D U^{\dagger} \tag{44}
\end{equation*}
$$

such that $D$ contains the eigenvalues. Here,

$$
\begin{equation*}
D=\operatorname{diag}\left(\lambda_{1}, \lambda_{2}, \lambda_{3}, \ldots, \lambda_{d}\right) \tag{45}
\end{equation*}
$$

and the eigenvectors are

$$
\begin{equation*}
U=\left[\left|\Psi_{1}\right\rangle,\left|\Psi_{2}\right\rangle,\left|\Psi_{3}\right\rangle, \ldots,\left|\Psi_{d}\right\rangle\right] . \tag{46}
\end{equation*}
$$

## Non-full-rank states II

- Assume that $\lambda_{d}=0$. Then,

$$
\begin{equation*}
\varrho^{\prime}=\varrho+\epsilon\left(\left|\Psi_{d}\right\rangle\left\langle\Psi_{d}\right|-\mathbb{1} / d\right) \tag{47}
\end{equation*}
$$

has a negative eigenvalue for any $\epsilon<0$. The Identity is needed to make the expression zero-trace.

- This is because the eigenvalues of this matrix are given by

$$
\begin{equation*}
D^{\prime}=\operatorname{diag}\left(\lambda_{1}-\epsilon / d, \lambda_{2}-\epsilon / d, \lambda_{3}-\epsilon / d, \ldots, \lambda_{d}+\epsilon(1-1 / d)\right) \tag{48}
\end{equation*}
$$

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## Two or more qubits: reduced states

- How can one see the state of a qubit, if it is the part of an entangled state?
- A reduced state of a bipartite system can be obtained after tracing out one of the subsystems. Let us consider a two-qubit system and write the density matrix in the basis $|00\rangle,|01\rangle,|10\rangle,|11\rangle$. Then, denote the elements of the density matrix by

$$
\begin{equation*}
\varrho_{i j, k l}, \tag{49}
\end{equation*}
$$

where $i, j, k, I=0,1$. In other words, it looks like

$$
\varrho=\left(\begin{array}{llll}
\varrho_{00,00} & \varrho_{00,01} & \varrho_{00,10} & \varrho_{00,11} \\
\varrho_{01,00} & \varrho_{01,01} & \varrho_{01,10} & \varrho_{01,11}  \tag{50}\\
\varrho_{10,00} & \varrho_{10,01} & \varrho_{10,10} & \varrho_{10,11} \\
\varrho_{11,00} & \varrho_{11,01} & \varrho_{11,10} & \varrho_{11,11}
\end{array}\right) .
$$

Thus, the size of the density matrix is $4 \times 4$.

## Two or more qubits: reduced states II

- To become familiar with bras and kets, one can even use the completeness relation

$$
\begin{equation*}
\text { Identity }=\sum_{i j}|i j\rangle\langle i j| . \tag{51}
\end{equation*}
$$

Then, one obtains

$$
\begin{equation*}
\text { Identity } \left.\times \varrho \times \text { Identity }=\sum_{i j k l}|i j\rangle(\langle i j| \varrho|k|\rangle\right)\langle k|, \tag{52}
\end{equation*}
$$

where the expression in the bracket is just the matrix element of the density matrix

$$
\begin{equation*}
\left.\varrho_{i j, k l}=\langle i j| \varrho|k\rangle\right\rangle . \tag{53}
\end{equation*}
$$

Hence, the density matrix can be written as

$$
\begin{equation*}
\varrho=\sum_{i j k l} \varrho_{i j, k \mid}|j\rangle\langle k| \mid . \tag{54}
\end{equation*}
$$

## Two or more qubits: reduced states III

- Then, tracing out the second subsystem gives the reduced state

$$
\begin{equation*}
\operatorname{Tr}_{2}(\varrho)=\varrho_{\mathrm{red}} \tag{55}
\end{equation*}
$$

which is given as

$$
\begin{equation*}
\varrho_{\mathrm{red}, \mathrm{ik}}=\sum_{m} \varrho_{i m, k m} . \tag{56}
\end{equation*}
$$

This is a $2 \times 2$ density matrix of a qubit. With this

$$
\begin{equation*}
\langle A \otimes \mathbb{1}\rangle_{\varrho}=\langle A\rangle_{\varrho_{\mathrm{red}}} . \tag{57}
\end{equation*}
$$

- Graphical representation: in the blockdiagonal representation, we sum the elements in the diagonal of the small matrices.
- Tracing out for pure states:

$$
\begin{equation*}
\operatorname{Tr}_{2}\left(\sum_{k} \alpha_{k}\left|\psi_{k}\right\rangle\left|\phi_{k}\right\rangle\right)=\sum_{k}\left|\alpha_{k}\right|^{2}\left|\psi_{k}\right\rangle\left\langle\psi_{k}\right| . \tag{58}
\end{equation*}
$$

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- Quantum relative entropy
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- K. Fidelity


## Purifications

- The pure state $\Psi_{A B}$ state is the purification of the mixed state $\varrho_{A}$ if

$$
\begin{equation*}
\operatorname{Tr}_{B}\left(\left|\Psi_{A B}\right\rangle\left\langle\Psi_{A B}\right|\right)=\varrho_{A} \tag{59}
\end{equation*}
$$

Note that $\left|\Psi_{A B}\right\rangle$ on subsystems $A$ and $B$, while $\varrho_{A}$ lives on subsystem $A$ only.

- Let us assume that a density matrix is defined as

$$
\begin{equation*}
\varrho_{A}=\sum_{k} p_{k}\left|\phi_{k}\right\rangle\left\langle\left.\phi_{k}\right|_{A} .\right. \tag{60}
\end{equation*}
$$

- Then, a purification can be a pure state

$$
\begin{equation*}
|\Psi\rangle_{A B}=\sum_{k} \sqrt{p_{k}}\left|\phi_{k}\right\rangle_{A} \otimes|k\rangle_{B} \tag{61}
\end{equation*}
$$

where $|k\rangle_{B}$ denotes an orthonormal basis of the subsystem $B$.

## Purifications II

- If $|\Psi\rangle_{A B}$ is a purification then

$$
|\Psi\rangle_{A B}^{\prime}=\mathbb{1}_{A} \otimes U_{B}|\Psi\rangle_{A B},
$$

is also a purification.

## Purifications III

- Purification of the eigendecomposition,

$$
\begin{equation*}
\varrho_{A}=\sum_{k} \lambda_{k}\left|\phi_{k}\right\rangle\left\langle\left.\phi_{k}\right|_{A} .\right. \tag{63}
\end{equation*}
$$

Then,

$$
\begin{equation*}
|\Psi\rangle_{A B}=\sum_{k} \sqrt{\lambda_{k}}\left|\phi_{k}\right\rangle_{A} \otimes|k\rangle_{B} . \tag{64}
\end{equation*}
$$

If $\varrho_{A}$ is full rank then the size of $B$ is the same of the size of $A$.

- In general, $B$ can also have a larger dimension that $A$.


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## Purity

- Defined as

$$
\begin{equation*}
\operatorname{Tr}\left(\varrho^{2}\right) \tag{65}
\end{equation*}
$$

- 1 for pure states.
- $1 / d$ for the completely mixed state.


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## Shannon entropy

- There is a source that outputs an integer number between 1 and $d$.
- The Shannon entropy is given as

$$
\begin{equation*}
H=-\sum_{k=1}^{d} p_{k} \log p_{k} \tag{66}
\end{equation*}
$$

## Shannon entropy II

- Properties
- Classical, not quantum.
- The source can have $d$ possible outputs with some probability.
- In information theory, the entropy of a random variable is the average level of "information", "surprise", or "uncertainty" inherent in the variable's possible outcomes (Wikipedia).
- There is a clear relation to compression of data. If the entropy is lower, one can compress the data to a smaller space.


## Shannon entropy III

- Further properties
- $H=0$ if $p_{1}=1$, all other $p_{k}=0 \cdot \vec{p}=(1,0,0,0, \ldots)$. The output is always the same. No information is provided.
- Comment: we can show that, using L'Hospitals rule,

$$
\begin{equation*}
\lim _{x \rightarrow 0}(x \log x)=\lim _{x \rightarrow 0} \frac{\log x}{1 / x}=\lim _{x \rightarrow 0} \frac{1 / x}{-1 / x^{2}}=-\lim _{x \rightarrow 0} x=0 \tag{67}
\end{equation*}
$$

- $H=\log d$ (maximal) if $p_{k}=\frac{1}{d} \cdot \vec{p}=\left(\frac{1}{d}, \frac{1}{d}, \frac{1}{d}, \frac{1}{d}, \ldots\right)$. All outputs are equally probable, a lot of information is provided.


## Von Neumann entropy

- Von Neumann entropy for a quantum state is defined as

$$
\begin{equation*}
S(\varrho)=-\operatorname{Tr}(\varrho \log \varrho) \equiv-\langle\log \varrho\rangle . \tag{68}
\end{equation*}
$$

- Note: matrix logarithm! It can be written with the eigenvalues of the density matrix as

$$
\begin{equation*}
S(\varrho)=-\sum_{k=1}^{d} \lambda_{k} \log _{2} \lambda_{k} \tag{69}
\end{equation*}
$$

## Von Neumann entropy II

- Properties
- Quantum. "Quantum version" of the Shannon entropy.
- For a pure state we have $\lambda_{k}=\{1,0,0, \ldots, 0\}$, and thus it is zero.
- Its maximal is for the completely mixed state for which $\lambda_{k}=\left\{\frac{1}{d}, \frac{1}{d}, \frac{1}{d}, \ldots, \frac{1}{d}\right\}$, and its value is $\log _{2} d$.
- Concave, i.e.,

$$
\begin{equation*}
S\left(p \varrho_{1}+(1-p) \varrho_{2}\right) \geq p S\left(\varrho_{1}\right)+(1-p) S\left(\varrho_{2}\right) . \tag{70}
\end{equation*}
$$

- Invariant under change of basis:

$$
\begin{equation*}
S(\varrho)=S\left(U_{\varrho} U^{\dagger}\right) \tag{71}
\end{equation*}
$$

## Von Neumann entropy III

- Further property
- Additive for independent systems.

$$
\begin{equation*}
S\left(\varrho_{1} \otimes \varrho_{2}\right)=-\operatorname{Tr}\left[\left(\varrho_{1} \otimes \varrho_{2}\right) \log \left(\varrho_{1} \otimes \varrho_{2}\right)\right] . \tag{72}
\end{equation*}
$$

Let us prove it. Let us first use

$$
\begin{equation*}
\log \left(\varrho_{1} \otimes \varrho_{2}\right)=\log \left(\varrho_{1}\right) \otimes \mathbb{1}+\mathbb{1} \otimes \log \left(\varrho_{2}\right) . \tag{73}
\end{equation*}
$$

We obtain

$$
S\left(\varrho_{1} \otimes \varrho_{2}\right)=-\operatorname{Tr}\left[\left(\varrho_{1} \otimes \varrho_{2}\right) \log \left(\varrho_{1}\right) \otimes \mathbb{1}\right]-\operatorname{Tr}\left[\left(\varrho_{1} \otimes \varrho_{2}\right) \mathbb{1} \otimes \log \left(\varrho_{2}\right)\right] .
$$

Let us now consider

$$
\begin{equation*}
(A \otimes B)(C \otimes D)=(A B) \otimes(C D) \tag{74}
\end{equation*}
$$

We obtain

$$
\begin{equation*}
S\left(\varrho_{1} \otimes \varrho_{2}\right)=-\operatorname{Tr}\left\{\left[\varrho_{1} \log \left(\varrho_{1}\right)\right] \otimes \varrho_{2}\right\}-\operatorname{Tr}\left\{\left[\varrho_{1} \otimes \varrho_{2} \log \left(\varrho_{2}\right)\right]\right\} . \tag{75}
\end{equation*}
$$

Finally, let us use

$$
\begin{equation*}
\operatorname{Tr}(A \otimes B)=\operatorname{Tr}(A) \operatorname{Tr}(B) \tag{76}
\end{equation*}
$$

We arrive at

$$
\begin{equation*}
S\left(\varrho_{1} \otimes \varrho_{2}\right)=-\operatorname{Tr}\left[\varrho_{1} \log \left(\varrho_{1}\right)\right]-\operatorname{Tr}\left[\varrho_{2} \log \left(\varrho_{2}\right)\right]=S\left(\varrho_{1}\right)+S\left(\varrho_{2}\right) . \tag{77}
\end{equation*}
$$

## Von Neumann entropy IV

- Further properties
- Strongly subadditive,

$$
\begin{equation*}
S\left(\varrho_{A B C}\right)+S\left(\varrho_{B}\right) \leq S\left(\varrho_{A B}\right)+S\left(\varrho_{B C}\right) . \tag{78}
\end{equation*}
$$

The matrices $\varrho_{B}, \varrho_{A B}$, etc. reduced states.

- Subadditive,

$$
\begin{equation*}
S\left(\varrho_{A C}\right) \leq S\left(\varrho_{A}\right)+S\left(\varrho_{C}\right) \equiv S\left(\varrho_{A} \otimes \varrho_{C}\right) . \tag{79}
\end{equation*}
$$

- Araki-Lieb inequality

$$
\begin{equation*}
\left|S\left(\varrho_{A}\right)-S\left(\varrho_{C}\right)\right| \leq S\left(\varrho_{A C}\right) . \tag{80}
\end{equation*}
$$

- Often used in condensed matter physics and field theory. See block entropy depending on the block size.


## Quantum relative entropy

- The relative entropy is given as

$$
\begin{equation*}
S(\varrho \| \sigma)=-\operatorname{Tr}[\varrho(\log \sigma-\log \varrho)]=-\operatorname{Tr}(\varrho \log \sigma)-S . \tag{81}
\end{equation*}
$$

- Properties
- $S(\varrho \| \sigma) \geq 0$.
- $S(\varrho \| \sigma)=0$ if and only if $\varrho=\sigma$.
- Not symmetric $S(\varrho \| \sigma) \neq S(\sigma \| \varrho)$.
- Sort of a distance between two quantum states.
- Invariant under simultaneous change of basis:

$$
S(\varrho \| \sigma)=S\left(U_{\varrho} U^{\dagger} \| U_{\sigma} U^{\dagger}\right)
$$

- $S\left(\varrho_{1} \otimes \varrho_{2} \| \sigma_{1} \otimes \sigma_{2}\right)=S\left(\varrho_{1} \| \sigma_{1}\right)+S\left(\varrho_{2} \| \sigma_{2}\right)$.


## Quantum relative entropy II

- Further properties
- For the relative entropy to the completely mixed state

$$
\begin{equation*}
\varrho_{\text {completely mixed }}=\mathbb{1} / d \tag{82}
\end{equation*}
$$

we have

$$
\begin{equation*}
S\left(\varrho \| \varrho_{\text {completely mixed }}\right)=\log (d)-S(\varrho) \tag{83}
\end{equation*}
$$

- Monotonicity under CP maps (completely positive maps = physical maps). $\varrho$ and $\sigma$ evolves under the same CP map. $S(\varrho \| \sigma)$ cannot increase.


## Linear entropy

- The linear entropy is defined as

$$
\begin{equation*}
S_{\operatorname{lin}}(\varrho)=1-\operatorname{Tr}\left(\varrho^{2}\right) \equiv\langle\mathbb{1}-\varrho\rangle . \tag{84}
\end{equation*}
$$

- It is often easier to obtain than the von Neumann entropy.
- Its relation to von Neumann entropy via the Mercator series is

$$
\begin{equation*}
-\langle\log \varrho\rangle=\langle\mathbb{1}-\varrho\rangle+\left\langle(\mathbb{1}-\varrho)^{2}\right\rangle / 2+\left\langle(\mathbb{1}-\varrho)^{3}\right\rangle / 3+\ldots \tag{85}
\end{equation*}
$$

This is based on expanding

$$
\begin{equation*}
\log (\mathbb{1}-(\mathbb{1}-\varrho)) \tag{86}
\end{equation*}
$$

using the Mercator series

$$
\begin{equation*}
\log (1+x)=x-x^{2} / 2+x^{3} / 3-+\ldots \tag{87}
\end{equation*}
$$

Note that

$$
\begin{equation*}
\mathbb{1}-\varrho \geq 0 \tag{88}
\end{equation*}
$$

Hence,

$$
\begin{equation*}
S \geq S_{\operatorname{lin}} \tag{89}
\end{equation*}
$$

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## Fidelity

How to measure the distance between quantum states?

- Pure states: overlap square, $\langle\psi \mid \phi\rangle^{2}$.
- $\langle\psi \mid \phi\rangle=0$ if and only if $|\psi\rangle=|\phi\rangle$.
- A pure state and a mixed state:

$$
\begin{equation*}
\operatorname{Tr}(|\Psi\rangle\langle\Psi| \varrho)=\langle\Psi| \varrho|\Psi\rangle . \tag{90}
\end{equation*}
$$

- Two mixed states: more difficult

$$
\begin{equation*}
F(\varrho, \sigma)=(\operatorname{Tr}(\sqrt{\sqrt{\varrho} \sigma \sqrt{\varrho}}))^{2} \tag{91}
\end{equation*}
$$

- $0 \leq F(\varrho, \sigma) \leq 1$.
- $F(\varrho, \sigma)=1$ if and only if $\varrho=\sigma$.
- $F(\varrho, \sigma)=0$ if $\varrho$ and $\sigma$ live on orthogonal subspaces.
- Symmetric $F(\varrho, \sigma)=F(\sigma, \varrho)$.
- Let us check consistency. If $\rho=|\Psi\rangle\langle\Psi|$ then $\sqrt{\varrho}=\varrho=|\Psi\rangle\langle\Psi|$. Then,

$$
\begin{equation*}
F(\varrho, \sigma)=\operatorname{Tr}(\sqrt{|\Psi\rangle\langle\Psi| \sigma|\Psi\rangle\langle\Psi|})^{2}=\langle\Psi| \sigma|\Psi\rangle \operatorname{Tr}(\sqrt{|\Psi\rangle\langle\Psi|})^{2}=\langle\Psi| \sigma|\Psi\rangle . \tag{92}
\end{equation*}
$$

Hence, we got back the formula for the simpler case.

## Fidelity

- Defining the Fidelity with a maximum over purifications

$$
\begin{equation*}
F(\varrho, \sigma)=\max _{\left|\Psi_{\sigma}\right\rangle}\left|\left\langle\Psi_{\varrho} \mid \Psi_{\sigma}\right\rangle\right|^{2} \tag{93}
\end{equation*}
$$

- $\left|\Psi_{\varrho}\right\rangle$ is a purification of $\varrho,\left|\Psi_{\sigma}\right\rangle$ is a purification of $\sigma$,

