Entanglement measures (How much is it entangled?) (Lecture of the Quantum Information class of the Master in Quantum Science and Technology)

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- Entanglement measures (How much is it entangled?)
 - Motivation
 - A. General quantum operation
 - B. Local operations and classical communication (LOCC)
 - C. Entanglement of formation
 - D. Concurrence
 - E. Entanglement of distillation
 - F. Bound entanglement
 - G. Requirements for entanglement measures
 - H. Negativity

Entanglement measures

 After detecting entanglement, we have to ask how entangled the state is.

It will turn out that entanglement is a resource.

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General quantum operation

The general quantum operation is defined as

$$\varrho' = \sum_{k} \mathsf{E}_{k} \varrho \mathsf{E}_{k}^{\dagger}$$

with

$$\sum_k E_k^{\dagger} E_k = 1.$$

- E_k are Kraus operators.
- Generalized measurements, POVM (positive operator-valued measure).
- Special case: von Neumann measurements, when E_k are pairwise orthogonal projectors.
- Naimark's dilation theorem: general operation= von Neumann measurement on system+ancilla.

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Local operations and classical communication (LOCC)

- LOCC are
 - local unitaries.
 - local von Neumann or POVM measurements,
 - local unitaries or measurements conditioned on measurement outcomes on the other party.
- Mathematical description of LOCC. Separable operations are a somewhat larger set, however, this set can easily be described.

$$\varrho' = \sum_{k} E_{k}^{(1)} \otimes E_{k}^{(2)} \varrho \left(E_{k}^{(1)} \otimes E_{k}^{(2)} \right)^{\dagger}$$

with

$$\sum_{k} \left(E_{k}^{(1)} \otimes E_{k}^{(2)} \right)^{\dagger} \left(E_{k}^{(1)} \otimes E_{k}^{(2)} \right) = 1.$$

Local operations and classical communication (LOCC) II

 Stochastic Local Operations and Classical Communication (SLOCC):

$$|\Psi\rangle' \leftarrow E_k^{(1)} \otimes E_k^{(2)} |\Psi\rangle$$

It happens with some probability, not deterministic.

- LOCC cannot create entanglement. Separable states remain separable under LOCC.
- LOCC can create correlations.

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Entropy of entanglement

The von Neumann entropy is defined as

$$S(\varrho) = -\text{Tr}(\varrho \log_2 \varrho).$$

It can be written with the eigenvalues of the density matrix as

$$S(\varrho) = -\sum_{k=1}^{d} \lambda_k \log_2 \lambda_k.$$

- For a pure state we have $\lambda_k = \{1, 0, 0, ..., 0\}$, and thus it is zero.
- Its maximal is for the completely mixed state for which $\lambda_k = \{\frac{1}{d}, \frac{1}{d}, \frac{1}{d}, ..., \frac{1}{d}\}$, and its value is $\log_2 d$.
- For a bipartite pure state, the entropy of entanglement is

$$E_E(|\Psi\rangle) = S(Tr_1(|\Psi\rangle\langle\Psi|)).$$

That is, it is the von Neumann entropy of the reduced state is an entaglement measure.

Entropy of entanglement II

- Comments
 - It is one for two-qubit singlet states.
 - It is zero for product states.
 - It is invariant under $U_1 \otimes U_2$.

Entanglement of formation

 For mixed states, the entanglement of formation is the convex roof of the von Neumann entropy of the reduced state.

$$E_F = \min_{|\Psi_k\rangle, p_k} \sum_k p_k E_E(|\Psi_k\rangle),$$

The optimization is over all decompositions of the state of the type

$$\varrho = \sum_{k} p_{k} |\Psi_{k}\rangle\langle\Psi_{k}|.$$

- *E_F* tells us, in the asymptotic limit, how many singlets we need to create the state.
- Is it easy to compute? No. For 2×2 systems, there is an explicit formula with the concurrence. For larger systems, there is not a general method.

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Entanglement of formation

- For two qubits, E_F can be calculated explicitly (Wootters, 1997).
- Special case: for pure states the concurrence is

$$C(|\Psi\rangle) = |\langle \Psi | \tilde{\Psi} \rangle| = 2|a_{11}a_{22} - a_{12}a_{21}|,$$

where

$$|\Psi
angle=\left(egin{array}{c} a_{11}\ a_{12}\ a_{21}\ a_{22} \end{array}
ight).$$

• It is related to the linear entropy of the reduced state.

$$C = \sqrt{2(1 - \text{Tr}(\rho_{\text{red}}^2)},\tag{1}$$

where

$$\rho_{\rm red} = \text{Tr}_2(|\Psi\rangle\langle\Psi|). \tag{2}$$

Entanglement of formation II

- Now we have to compute E_F from C.
- We also nee that

$$\epsilon(c) = H_2\left(\frac{1+\sqrt{1-c^2}}{2}\right).$$

Here

$$H_2 = -x \log_2 x - (1-x) \log_2 (1-x).$$

• Then, E_F can be obtained as

$$E_F(\varrho) = \epsilon(C(\varrho)).$$

Entanglement of formation III

For mixed states, the concurrence is defined as

$$C(\varrho) = \max(0, \lambda_1 - \lambda_2 - \lambda_3 - \lambda_4),$$

where λ_k 's are, in a decreasing order, the eigenvalues of

$$R=\sqrt{\sqrt{\varrho}\tilde{\varrho}\sqrt{\varrho}},$$

and

$$\tilde{\varrho} = (\sigma_y \otimes \sigma_y) \varrho^* (\sigma_y \otimes \sigma_y).$$

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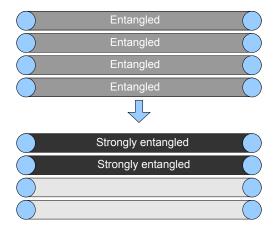
Entanglement of distillation

 E_D tells us, how many singlets we can obtain from the state with LOCC. In general,

$$E_F \geq E_D$$
.

 Note that local operation and classical communication means that we have several copies and we can act on the copies locally.

Entanglement of distillation II



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Bound entanglement

- There are states that need entangled particles to be created, but singlets cannot be distilled from them.
- All PPT entangled states are like that. (That is, all entangled states that are not detected by the Peres-Horodecki criterion.)

Bound entanglement II

 Next, we will prove this. First we show that PPT state remain PPT under LOCC. Under LOCC we have

$$\varrho' = \sum_{k} E_k^{(1)} \otimes E_k^{(2)} \varrho \left(E_k^{(1)} \otimes E_k^{(2)} \right)^{\dagger}$$

We also have

$$(\varrho')^{T2} = \sum_{k} E_{k}^{(1)} \otimes ((E_{k}^{(2)})^{\dagger})^{T} \varrho^{T2} (E_{k}^{(1)})^{\dagger} \otimes (E_{k}^{(2)})^{T}$$

Here we used that $(AB)^T = B^T A^T$ and $A^{\dagger} = (A^*)^T$.

• We can see that if $\varrho^{T2} \ge 0$ then $(\varrho')^{T2} \ge 0$. Thus the PPT states remain PPT under LOCC.

R., P., M., and K. Horodecki, Rev. Mod. Phys. 81, 865 (2009). (Click on the link above, see "G. Bound entanglement - when distillability fails" on page 44.)

Bound entanglement III

• Let us again remember the flip operator

$$F|k\rangle|l\rangle = |l\rangle|k\rangle$$

It has eigenvalues ±1.

• The maximally entangled state

$$|\Psi_{\rm me}\rangle = \frac{1}{\sqrt{d}} \sum_{k=1}^{d} |k\rangle |k\rangle.$$

Bound entanglement IV

We can show that

$$\begin{split} |\Psi_{\rm me}\rangle\langle\Psi_{\rm me}| &= \frac{1}{d}\sum_{k,l}^{d}|k\rangle\langle l|\otimes|k\rangle\langle l|,\\ |\Psi_{\rm me}\rangle\langle\Psi_{\rm me}|^{T1} &= \frac{1}{d}\sum_{k,l}^{d}|k\rangle\langle l|\otimes|l\rangle\langle k|\equiv\frac{F}{d}. \end{split}$$

 Now we show that PPT states have a small overlap with the maximally entangled state. For PPT states, the fidelity with respect to the maximally entangled state is

$$\operatorname{Tr}(|\Psi_{\mathrm{me}}\rangle\langle\Psi_{\mathrm{me}}|\varrho) = \operatorname{Tr}(|\Psi_{\mathrm{me}}\rangle\langle\Psi_{\mathrm{me}}|^{T_1}\varrho^{T_1}) = \frac{1}{d}\operatorname{Tr}(F\varrho^{T_1}) \leq \frac{1}{d},$$

since $\varrho^{T1} \ge 0$ and F has ± 1 eigenvalues.

Bound entanglement IV

- Thus, PPT states have a small fidelity with respect to the maximally entangled state. Even LOCC operations cannot increase this.
- A simple product state can reach 1/d

$$\operatorname{Tr}(|\Psi_{\mathrm{me}}\rangle\langle\Psi_{\mathrm{me}}||11\rangle\langle11|) = \frac{1}{d}.$$

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Requirements for entanglement measures

- To each density matrix it assigns a nonnegative number. Typically, the maximally entangled state has log d.
- ② $E(\varrho) = 0$ for separable states.
- 3 E does not increase on average under LOCC.

$$E(\varrho) \le \sum_{k} p_{k} E\left(\frac{A_{k} \varrho A_{k}^{\dagger}}{\operatorname{Tr}(A_{k} \varrho A_{k}^{\dagger})}\right). \tag{3}$$

- For pure states, it has the same value as the entangement entropy.
 - Entanglement monotone: 1,2,3.
- Entanglement mesure: 1,2,4 and does not increase under deterministic LOCC, i.e.,

$$E(\varrho') \le E(\varrho); \quad \varrho' = \sum_{k} A_k \varrho A_k^{\dagger} \quad \text{(POVM)}.$$
 (4)

M. B. Plenio and S. Virmani, eprint arXiv:quant-ph/0504163 (2005).

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Trace norm

Let us consider the singular decomposition of a matrix

$$A = U \Sigma V^{\dagger}$$
,

where

$$\Sigma = \operatorname{diag}(\sigma_1, \sigma_2, \sigma_3, ..., \sigma_d)$$

Then the trace norm is

and $\sigma_k > 0$.

The Hilbert-Schmidt norm is

$$\sigma_2 = \sigma_4$$

$$||A||_1 = \operatorname{Tr}\left(\sqrt{AA^{\dagger}}\right) = \sum_k \sigma_k.$$

(5)

(6)

(7)

$$||A||_2 = T$$

$$||A||_2 = \operatorname{Tr}(AA^{\dagger}) = \sum_k \sigma_k^2.$$

Negativity

Example for a monotone: negativity

$$N(\varrho) = \frac{\|\varrho^{\mathrm{T1}}\| - 1}{2}.$$

Trace norm=sum of singular values.

• For Hermitian matrices, it is the same as sum of eigenvalues.

$$N(\varrho) = \frac{\sum_{k} |\lambda_{k}| - 1}{2}.$$

• Note that $\sum_k \lambda_k = 1$. Then, assume that the first M eigenvalues are negative, the rest is positive. We get

$$N(\varrho) = \frac{\sum_{k=1}^{M} -\lambda_k + \sum_{k=M+1}^{d} \lambda_k - \sum_k \lambda_k}{2}.$$

Negativity II

Hence,

$$N(\varrho) = \sum_{k=1}^{M} |\lambda_k|.$$

That is, the absolute value of the sum of the negative eigenvalues of the partial transpose.

- Clearly, it is zero for PPT states. Thus, it is zero for all separable states.
- Not as meaningful as the Entanglement of Formation, but can be calculated on any system sizes.
- It fulfills certain conditions on how it changes under LOCC. It does not increase under deterministic LOCC.